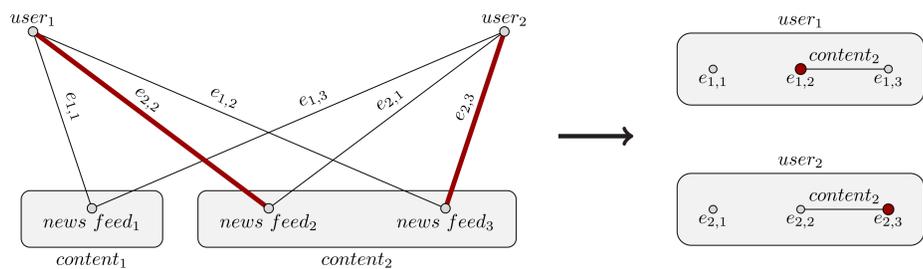


EFFICIENT LEARNING BY IMPLICIT EXPLORATION IN BANDIT PROBLEMS WITH SIDE OBSERVATIONS

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MOTIVATION - SEQUENTIAL NEWS RECOMMENDATION



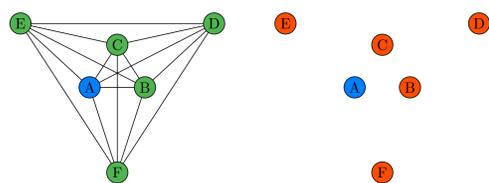
- Select a matching covering users
- Obtain rewards of selected edges
- Observe additional rewards

APPLICATIONS

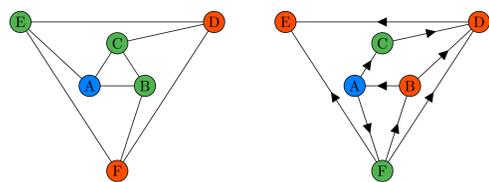
- **Packet routing in computer networks**
 - Typical feedback: delays of our own packets
 - Side observations: other delays in network
- **Web advertising – displaying one ad**
 - Case 1: symmetric interrelations
 - ◊ Example: similar preferences for similar cars
 - ◊ Typical feedback: reward for displayed ad
 - ◊ Side observations for similar cars
 - ◊ Model for observations: undirected graph
 - Case 2: asymmetric interrelations
 - ◊ Example: electronics (interest in camera means interest in accessories, not vice versa)
 - ◊ Typical feedback: reward for displayed ad
 - ◊ Side observations for dependent products
 - ◊ Model for observations: directed graph

EXAMPLES OF GRAPH STRUCTURES

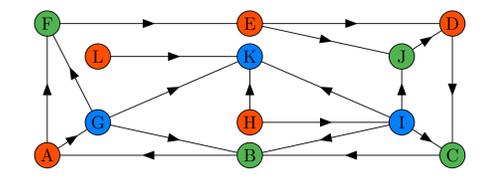
Side observations can be modeled as a graph
Full information and bandit setting – simple action



Undirected and directed case – simple action



Directed combinatorial case – complex action



LEARNING SETTING

In every round $t = 1, 2, \dots, T$:

- **Environment**
 - Privately assigns vector ℓ_t of losses to actions
 - Generates an observation graph
 - ◊ Undirected / Directed
 - ◊ Disclosed / Not disclosed
- **Learner**
 - Plays action $V_t \in \mathcal{S} \subseteq \{0, 1\}^N$
 - ◊ Each action $v \in \mathcal{S}$ satisfies $\|v\|_1 \leq m$
 - ◊ I.e. action consists of playing at most m nodes
 - ◊ Case $m = 1$: we denote $I_t \in [N]$ a node played
 - Obtain loss $V_t^\top \ell_t$ corresponding to nodes played
 - Observe losses of neighbors of played nodes
 - ◊ Graph disclosed

Performance measure: total expected regret

$$R_T = \max_{v \in \mathcal{S}} \mathbb{E} \left[\sum_{t=1}^T (V_t - v)^\top \ell_t \right]$$

IMPLICIT EXPLORATION

Usual approach to exploration:

- Bias sampling distribution as $\tilde{p}_t = (1 - \gamma)p_t + \gamma\mu$
 - Needs to know graph structure
 - Constructing a good μ is expensive
- Construct unbiased loss estimates

$$\hat{\ell}_{t,i} = \frac{\ell_{t,i}}{o_{t,i}} \mathbb{1}\{\ell_{t,i} \text{ is observed}\}$$
- $o_{t,i}$ – probability of observing $\ell_{t,i}$

Our new approach:

- Do not touch sampling distribution
- Construct **optimistically biased** loss estimates

$$\hat{\ell}_{t,i} = \frac{\ell_{t,i}}{o_{t,i} + \gamma} \mathbb{1}\{\ell_{t,i} \text{ is observed}\} \quad \text{s.t.} \quad \mathbb{E}[\hat{\ell}_{t,i}] \leq \ell_{t,i}$$
- Encourages exploration by optimism
- Does not require knowledge of observation graph
- Cheaper than computing μ

EXP3-IX ALGORITHM

- Compute weights using loss estimates $\hat{\ell}_{t,i}$.

$$w_{t,i} = \exp \left(-\eta \sum_{s=1}^{t-1} \hat{\ell}_{s,i} \right)$$

- Play action I_t according to probability distribution

$$\mathbb{P}(I_t = i) = p_{t,i} = \frac{w_{t,i}}{W_t} = \frac{w_{t,i}}{\sum_{j=1}^N w_{t,j}}$$

- Compute loss estimates (using observability graph)

$$\hat{\ell}_{t,i} = \frac{\ell_{t,i}}{o_{t,i} + \gamma} \mathbb{1}\{\ell_{t,i} \text{ is observed}\}$$

FPL-IX ALGORITHM

- Draw perturbation $Z_{t,i} \sim \text{Exp}(1)$ for all $i \in [N]$
- Play “the best” action V_t according to total loss estimate \hat{L}_{t-1} and perturbation Z_t

$$V_t = \arg \min_{v \in \mathcal{S}} v^\top (\eta_t \hat{L}_{t-1} - Z_t)$$

- Compute loss estimates

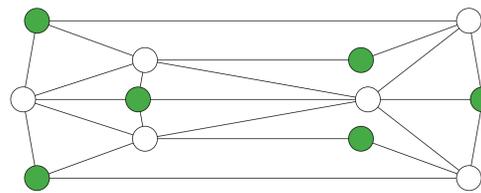
$$\hat{\ell}_{t,i} = \ell_{t,i} K_{t,i} \mathbb{1}\{\ell_{t,i} \text{ is observed}\}$$

- $K_{t,i}$: geometric random variable with

$$\mathbb{E}[K_{t,i}] = \frac{1}{o_{t,i} + (1 - o_{t,i})\gamma}$$

INDEPENDENCE SET

- Nodes of independence set are not connected
- α - size of the largest independence set



Independence set of size 6

MAIN RESULTS

Regret bound of Exp3-IX

$$R_T = \tilde{O} \left(\sqrt{\sum_{t=1}^T \alpha_t} \right) = \tilde{O}(\sqrt{\alpha T})$$

$\bar{\alpha}$ - average independence number of observation graph

Regret bound of FPL-IX

$$R_T = \tilde{O} \left(m^{3/2} \sqrt{\sum_{t=1}^T \alpha_t} \right) = \tilde{O}(m^{3/2} \sqrt{\alpha T})$$

RELATED WORK

- **Undirected case – simple action ($m = 1$)**
 - ELP (Mannor, Shamir)
 - ◊ Graph disclosed before action
 - ◊ Need to compute linear program for mixing
 - ◊ Regret bound of order $\tilde{O}(\sqrt{cT})$
 - **Exp3-IX**
 - ◊ Graph disclosed after action
 - ◊ Computationally efficient
 - ◊ Regret bound of order $\tilde{O}(\sqrt{\alpha T})$
- **Directed case - complex action ($m > 1$)**
 - **FPL-IX**
 - ◊ Graph disclosed after action
 - ◊ Computationally efficient
 - ◊ Regret bound of order $\tilde{O}(m^{3/2} \sqrt{\alpha T})$

ANALYSIS

Analysis of Exp3 algorithms in general - tracking evolution of $\log(W_{t+1}/W_t)$

$$\mathbb{E} \left[\underbrace{\sum_{t=1}^T \sum_{i=1}^N p_{t,i} \hat{\ell}_{t,i}}_A \right] - \mathbb{E} \left[\underbrace{\sum_{t=1}^T \sum_{k=1}^N \hat{\ell}_{t,k}}_B \right] \leq \mathbb{E} \left[\frac{\log N}{\eta} \right] + \mathbb{E} \left[\underbrace{\frac{\eta}{2} \sum_{t=1}^T \sum_{i=1}^N p_{t,i} (\hat{\ell}_{t,i})^2}_C \right]$$

Lower bound of A (using definition of loss estimates)

$$\mathbb{E} \left[\sum_{t=1}^T \sum_{i=1}^N p_{t,i} \hat{\ell}_{t,i} \right] \geq \mathbb{E} \left[\sum_{t=1}^T \sum_{i=1}^N p_{t,i} \ell_{t,i} \right] - \mathbb{E} \left[\gamma \sum_{t=1}^T Q_t \right]$$

Lower bound of B (optimistic loss estimates: $\mathbb{E}[\hat{\ell}] < \mathbb{E}[\ell]$)

$$-\mathbb{E} \left[\sum_{t=1}^T \hat{\ell}_{t,k} \right] \geq -\mathbb{E} \left[\sum_{t=1}^T \ell_{t,k} \right]$$

Upper bound of C (using definition of loss estimates)

$$\mathbb{E} \left[\frac{\eta}{2} \sum_{t=1}^T \sum_{i=1}^N p_{t,i} (\hat{\ell}_{t,i})^2 \right] \leq \mathbb{E} \left[\frac{\eta}{2} \sum_{t=1}^T Q_t \right]$$

Together we have

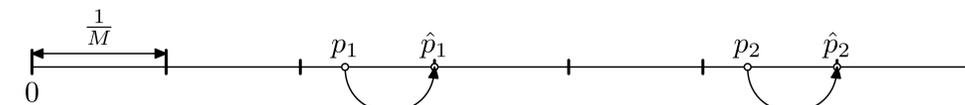
$$R_T \leq \frac{\log N}{\eta} + \left(\frac{\eta}{2} + \gamma \right) \sum_{t=1}^T \mathbb{E}[Q_t]$$

$$Q_t = \sum_{i=1}^N \frac{p_{t,i}}{o_{t,i} + \gamma}$$

Lemma 1. Let G be a directed graph, with $V = \{1, \dots, N\}$. Let d_i^- be the indegree of the node i and $\alpha = \alpha(G)$ be the independence number of G . Then

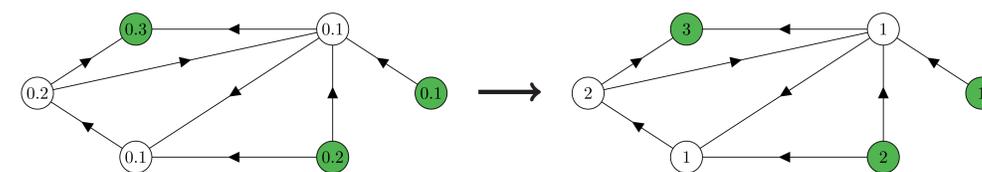
$$\sum_{i=1}^N \frac{1}{1 + d_i^-} \leq 2\alpha \log \left(1 + \frac{N}{\alpha} \right).$$

Step 1 of applying Lemma 1 to upper bound Q_t - Discretization



$$Q_t = \sum_{i=1}^N \frac{p_{t,i}}{o_{t,i} + \gamma} = \sum_{i=1}^N \frac{p_{t,i}}{p_{t,i} + \sum_{j \in N_i^-} p_{t,j} + \gamma} \leq \sum_{i=1}^N \frac{\hat{p}_{t,i}}{\hat{p}_{t,i} + \sum_{j \in N_i^-} \hat{p}_{t,j}} + 2 \quad \text{for } M = \lceil N^2/\gamma \rceil$$

Step 2 of applying Lemma 1 to upper bound Q_t - Construction of a “clique graph”



$$\sum_{i=1}^N \frac{\hat{p}_{t,i}}{\hat{p}_{t,i} + \sum_{j \in N_i^-} \hat{p}_{t,j}} = \sum_{i=1}^N \frac{M \hat{p}_{t,i}}{M \hat{p}_{t,i} + \sum_{j \in N_i^-} M \hat{p}_{t,j}} = \sum_{i=1}^N \sum_{k \in C_i} \frac{1}{1 + d_k^-} \leq 2\alpha \log \left(1 + \frac{M + N}{\alpha} \right)$$