

# Cheap Bandits

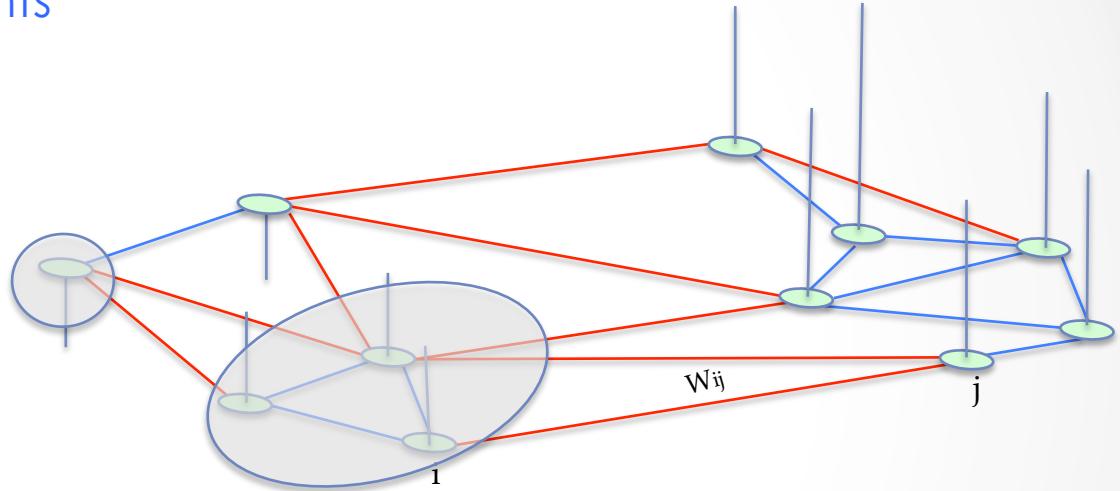
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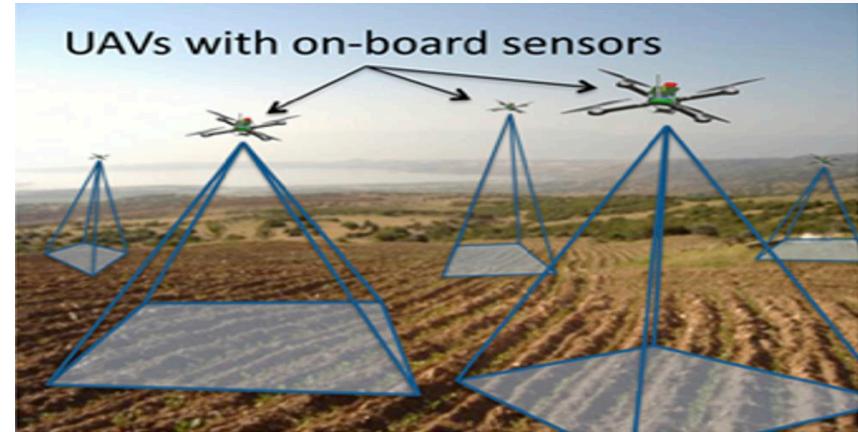
# Problem setting

- Undirected Graph:  $G=(V,E,W)$ 
  - N Nodes,  $W=\{w_{ij}\}$ : Weights
- Signal on Graph
  - Reward Function  
 $f : V \rightarrow \mathbb{R}$
  - Smooth Function
- Locate maxima
  - $$u^* = \arg \max_{u \in V} f(u)$$
- Actions:
  - Noisy Cluster Averages; Differentiated Costs
- Goal: In min Time ( $T \ll N$ ) locate  $u^*$ ; Min Cost?



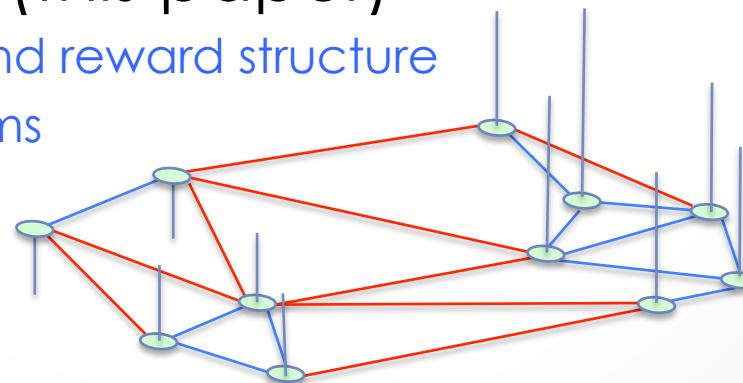
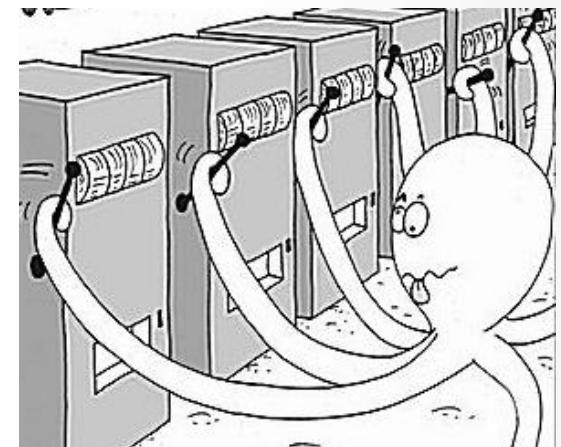
# Application Scenarios

- Surveillance/Geography
  - Forest Cover Dataset: labeled samples on  $30m^2$  region
  - **Nodes**: Regions of forest; **Edge weights**: feature similarity;
  - **Rewards**: Density of species. Locate highest density.
  - **Actions**: Zoom-in to a node (high cost); Zoom-out (low cost).
- Sensor networks:
- Radar search:
- Online advertisements:



# Bandit Setting

- N-arm Bandit [Robbins'72, Lai-Robbins85]
  - N Independent Rewards/arms
    - Each arm  $\sim$  action
  - N-nodes  $\sim$  no coupling between nodes
  - Need  $T \gg N$ .
    - Multiple looks per node
- We want  $T \ll N$  (this paper)
  - Exploit graph and reward structure
  - Very large # arms



# Reward is Linear and Smooth

- Linear Reward
  - Fourier decomposition

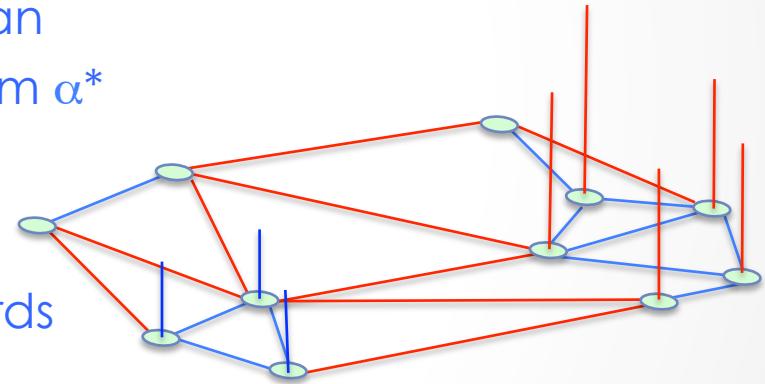
$$f = Q\alpha^*$$

- $Q$ : Eigenvectors of the graph Laplacian
- Linearly Param Bandit: unknown param  $\alpha^*$

- Smooth Reward

- Neighboring nodes have similar rewards

$$(u, v) \in E \implies f(u) \approx f(v)$$



$$\|\mathcal{L}f\|_2^2 = \sum_{u,v} w_{uv} (f(u) - f(v))^2 \leq c$$

[Valko et. al. ICML'14]

# Actions: Sample Node or Group

- Actions consists of subset of simplex:

- Sample a node,  $u$ :

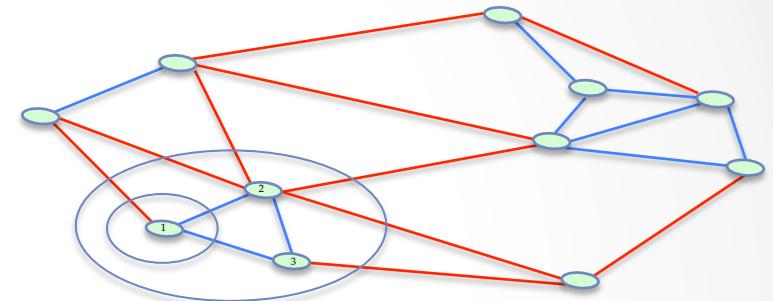
$$s(v) = \delta(u - v)$$

- Sample a group of nodes  $A \subset V$

$$s(v) = \frac{1}{|A|} \sum_{u \in A} \delta(u - v)$$

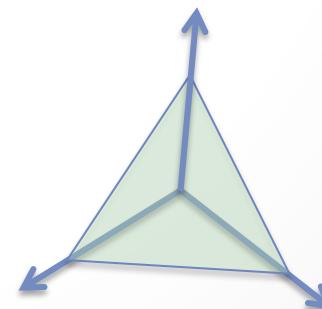
- General
    - Any Probability Mass Function

$$\mathcal{S} = \Delta^N$$



$$[1 \ 0 \ 0 \ 0 \ \dots \ 0]$$

$$\left[ \frac{1}{3} \ \frac{1}{3} \ \frac{1}{3} \ 0 \ \dots \ 0 \right]$$



# Cost of Actions

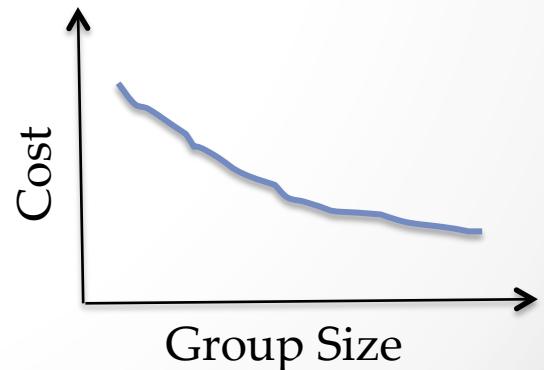
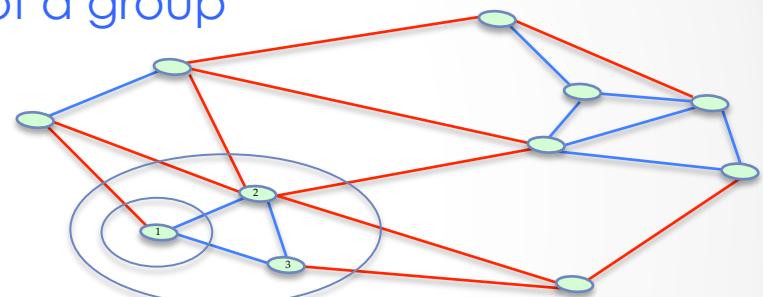
- Cost of actions:

- Costly: Zoom-in to observe a particular node
- Cheap: Zoom-out to observe average of a group

- Cost Model

$$C(s) = \sum_{(u,v) \in E} (s(u) - s(v))^2 = \|\mathcal{L}s\|_2^2$$

- Why this model?
  - Larger the group size smaller the cost
  - Probing Nodes has high cost
  - In Fourier domain: Energy of  $s$



# Regret and Cost

- Policy( $\pi$ ): In round  $t$ , pick an action  $s_t$

- Observe reward

$$r_t(s_t) = \langle s_t, f \rangle + \epsilon_t = \sum_u s_t(u) f(u) + \epsilon_t$$

- Cumulative Regret

$$R_T(\pi) = T f(u^*) - E \left[ \sum_{t=1}^T r_t(s_t) \right]$$

- Cumulative Cost

$$C_T(\pi) = \sum_{t=1}^T C(s_t)$$

# Objective: Cost vs Regret

- Minimize Cost subject to ‘optimal’ Regret

$$\min_{\pi, S} C_T(\pi)$$

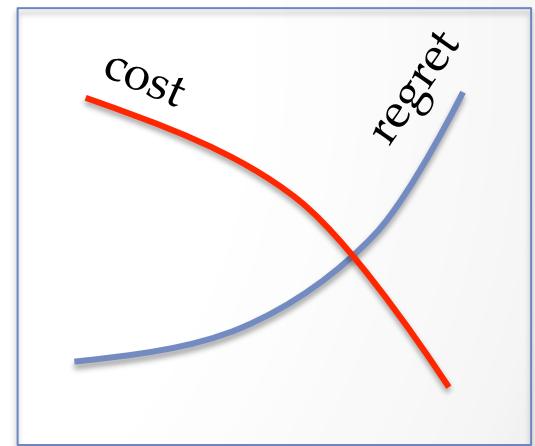
$$\text{subject to } R_T(\pi) \leq R_T^*$$

- Best admissible policies

$$R_T^* = \min_{\pi, S} R_T(\pi)$$

- Conflicting goals:

- Node actions give better estimates, but costly
- Group actions give poor estimates, but cheaper



Optimal Regret with lower cost

# What is a good Regret Constraint? Lower Bound

- No smoothness constraint ( $c \rightarrow \infty$ )

- Finite set of actions

$$R_T(\cdot) = \Omega(\sqrt{NT}) \quad (\text{Chu et. al. AISTATS'11})$$

- Smooth Functions (This paper)

**Proposition:** For Smooth function on graphs with effective dimension  $d$

$$R_T(\cdot) = \Omega(\sqrt{dT}) \quad \text{where} \quad d \ll N$$

- Effective Dimension [Valko et.al. ICML'14]

$$d = \max \left\{ i \mid \lambda_i(i-1) \leq \frac{T}{\log(T+1)} \right\}$$

# Intuition: Lower Bound

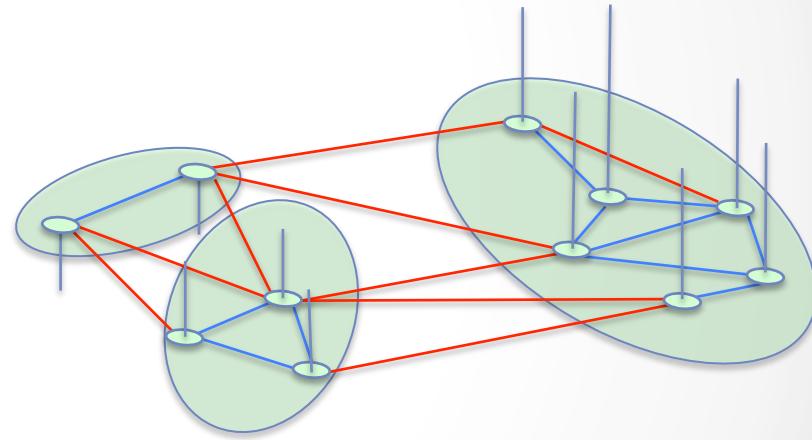
- Effective Dimension related to Graph Clusters

- d clusters

- # Disconnected clusters or
  - # sparse clusters

$$d = \max \left\{ i \mid \lambda_i(i-1) \leq \frac{T}{\log(T+1)} \right\}$$

Need to sample at least one node per cluster



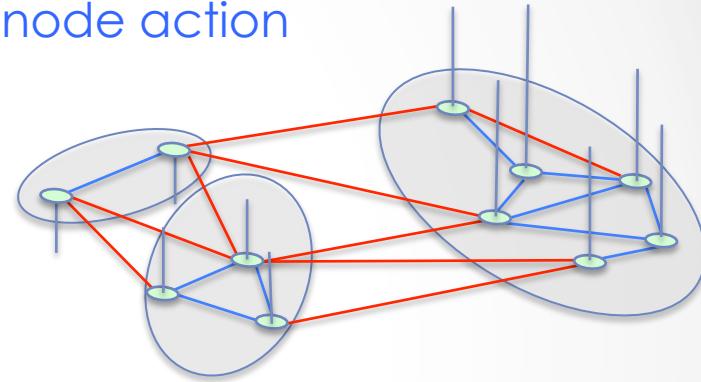
$$\min_{\pi, \mathcal{S}} C_T(\pi)$$

subject to  $R_T(\pi) \leq \mathcal{O}(\sqrt{dT})$

# Key Intuition: Locally Smooth Rewards

- Smoothness condition implies local smoothness
  - Group actions are good approximation to node action

$$u \in A \implies f(u) \sim \frac{1}{|A|} \sum_{v \in A} f(v) + \text{const}$$

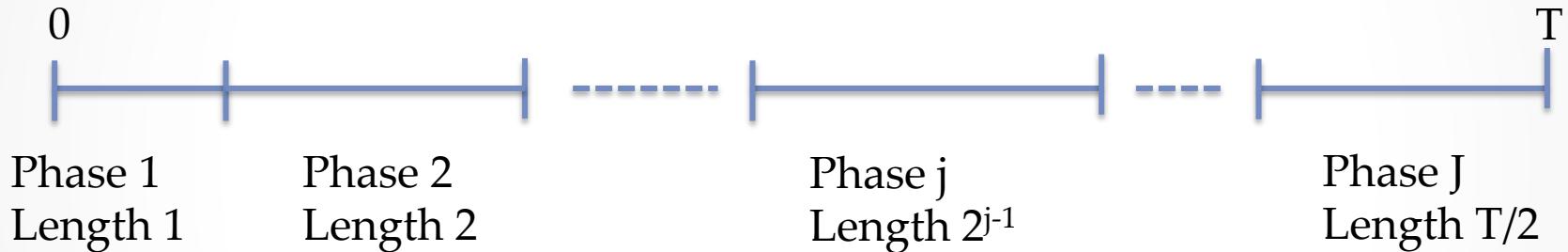


**Proposition:** Let  $\mathbf{f}$  be a smooth function on a graph with effective dimension  $d$ . Then,

$$|\mathbf{f}(i) - \frac{1}{\mathcal{N}_i} \sum_{j \in \mathcal{N}_i} \mathbf{f}(j)| \leq \frac{c'd}{\lambda_{d+1}}$$

# CheapUCB: Algorithm

- Inspired by SpectralUCB Algorithm [Valko et. al. ICML14]
- SpectralUCB uses only node actions, cannot control cost
- CheapUCB uses both node actions and group actions



- **Phases:** Split the  $T$  into  $J = \lceil \log T \rceil$  phases
- **Length:** Phase  $j=1,2,\dots,J$  is of  $2^{j-1}$  rounds
- **Select action:** In phase  $j$  select groups of size  $J-j+1$  optimistically using UCB

**Zoom-in slowly using progressively costly actions**

# Algorithm Performance

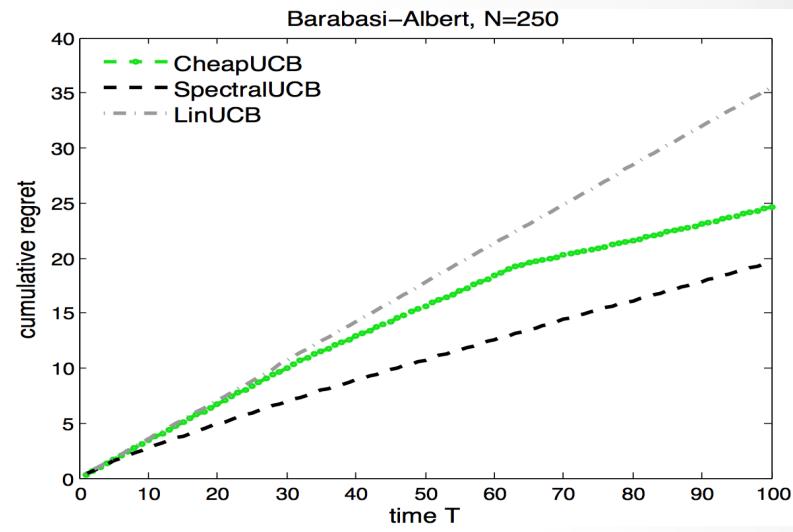
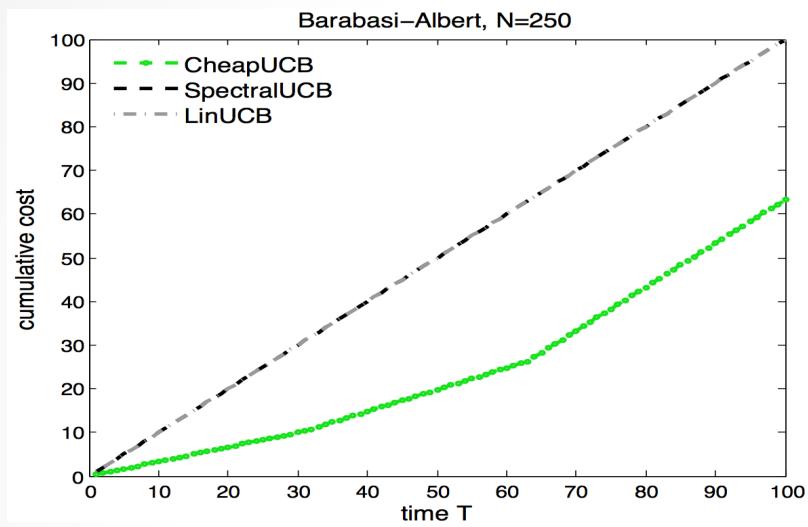
Algorithm	Regret bound	Cost
SpectralUCB (ICML'14)	$\mathcal{O}(d\sqrt{T})$	T
CheapUCB (This paper)	$\mathcal{O}(d\sqrt{T})$	$\frac{3}{4} T$

CheapUCB provides good regret guarantee and also provides  $O(T)$  cost saving

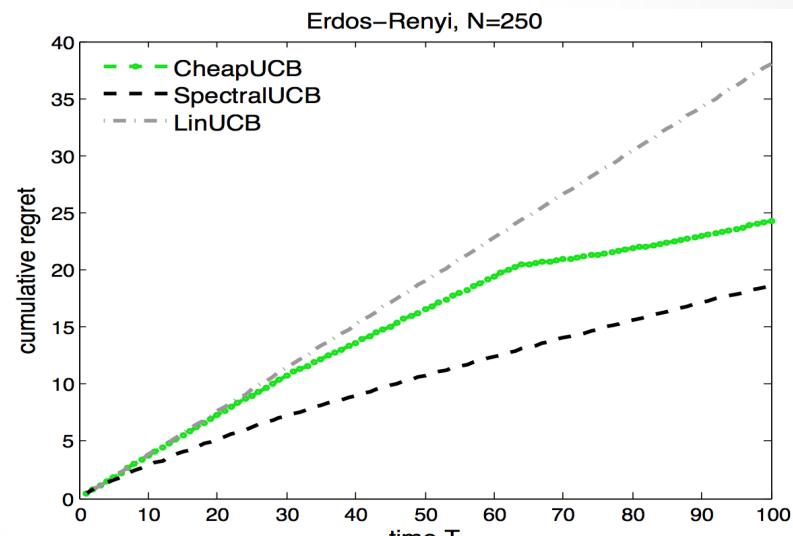
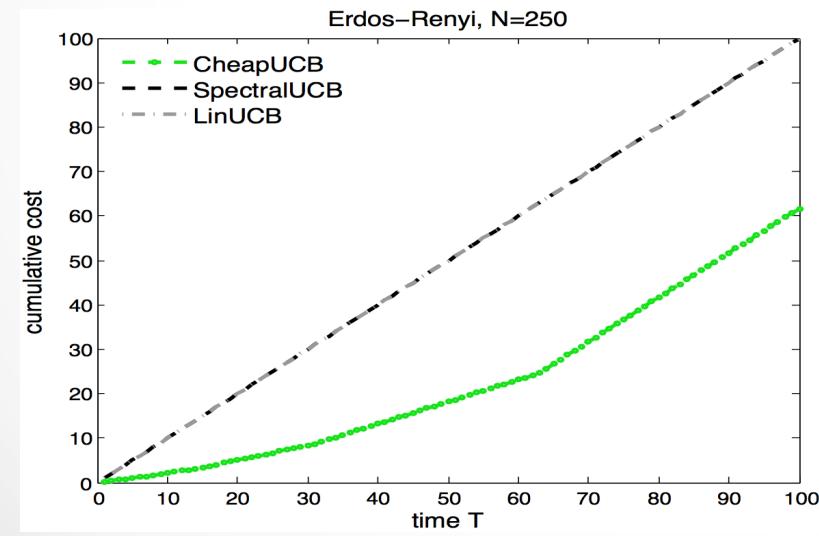
CheapUCB achieves at least 25% reduction in cost!!

# Network Experiments

Barabasi-Albert

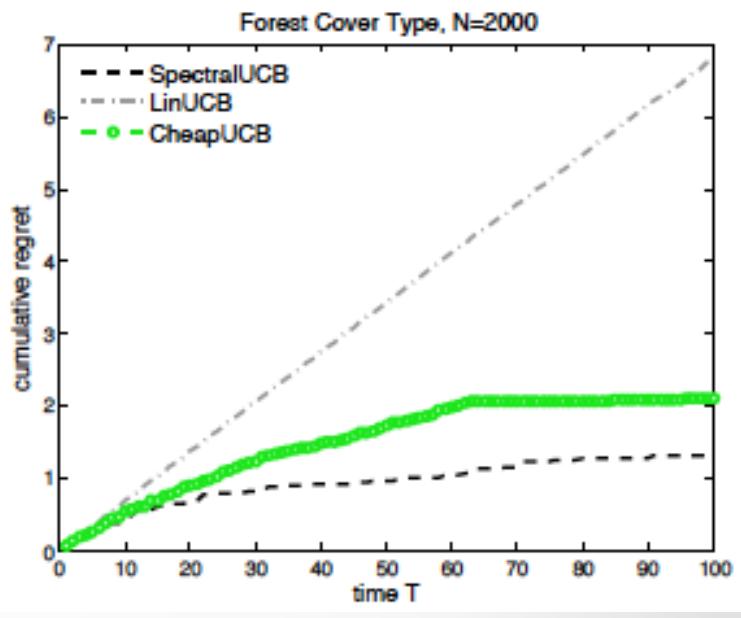
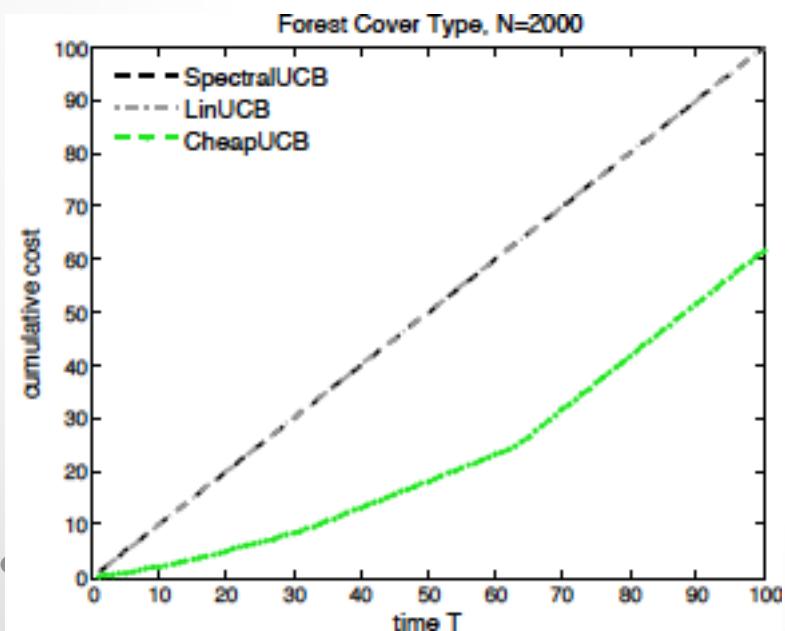
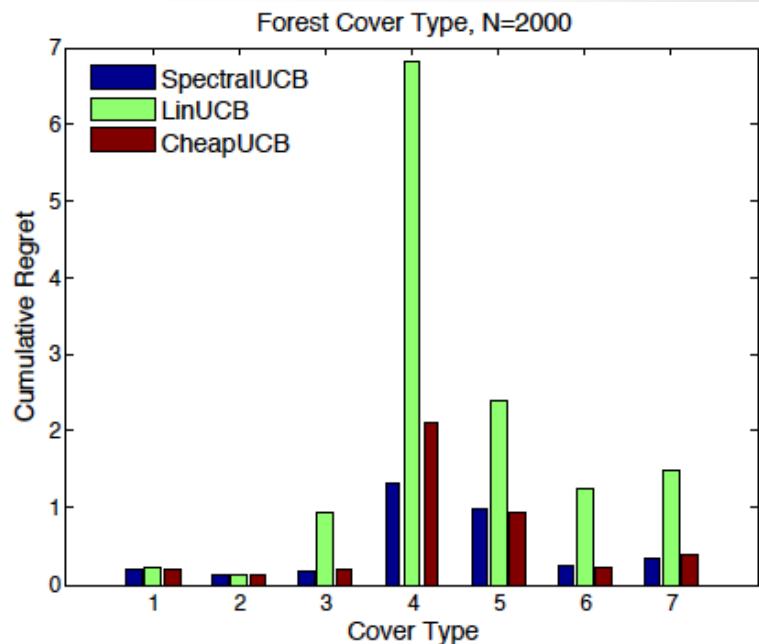


Erdos-Renyi



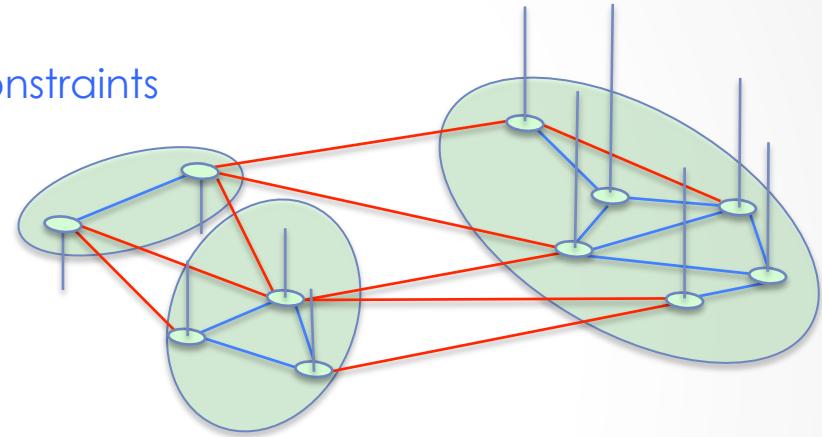
# Forest Cover Dataset

- 50000 Samples; 7 Species
- 30m<sup>2</sup> regions; 2000 clusters
- Nodes: regions; Edges: Feature similarity
  - Connect K-NN
- Reward: Density of Desired Species
  - Continuous Classifier Output



# Conclusions

- Cheap Bandit Formulation
  - Optima of Smooth signals on graphs
  - Minimize cost under optimal regret constraints
- Probes/Actions
  - Actions: Sample a node or a group
  - Cost of actions
- Effective Dimension governs regret
  - Time  $\ll N$ , depends on statistical dimension
- Expand actions beyond node actions to reduce cost
  - CheapUCB algorithm
  - Reduces cost by at least by 25%



# **Thank You!!**

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