

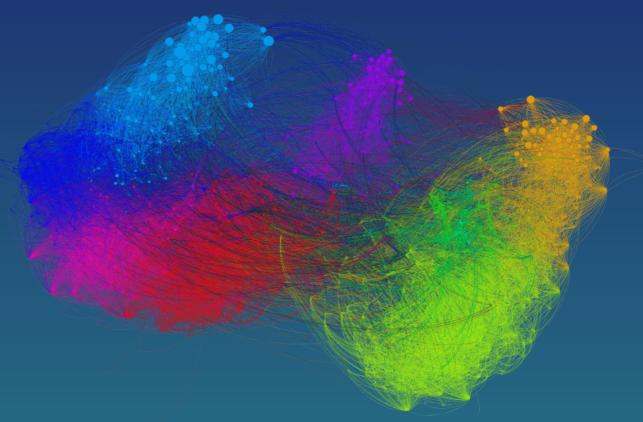
Graphs in Machine Learning

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Partially based on material by: Tomáš Kocák, Nikhil Srivastava, Yiannis Koutis, Joshua Batson, Daniel Spielman



December 12, 2016 MVA 2016/2017

LAST LECTURE



- Scaling harmonic functions to millions of samples
- Graph Sparsification
- Spectral Sparsification

THIS LECTURE



- Graph bandits
 - Spectral bandits
 - Observability graphs
 - Side information
 - ► Influence Maximization

PREVIOUS LAB SESSION

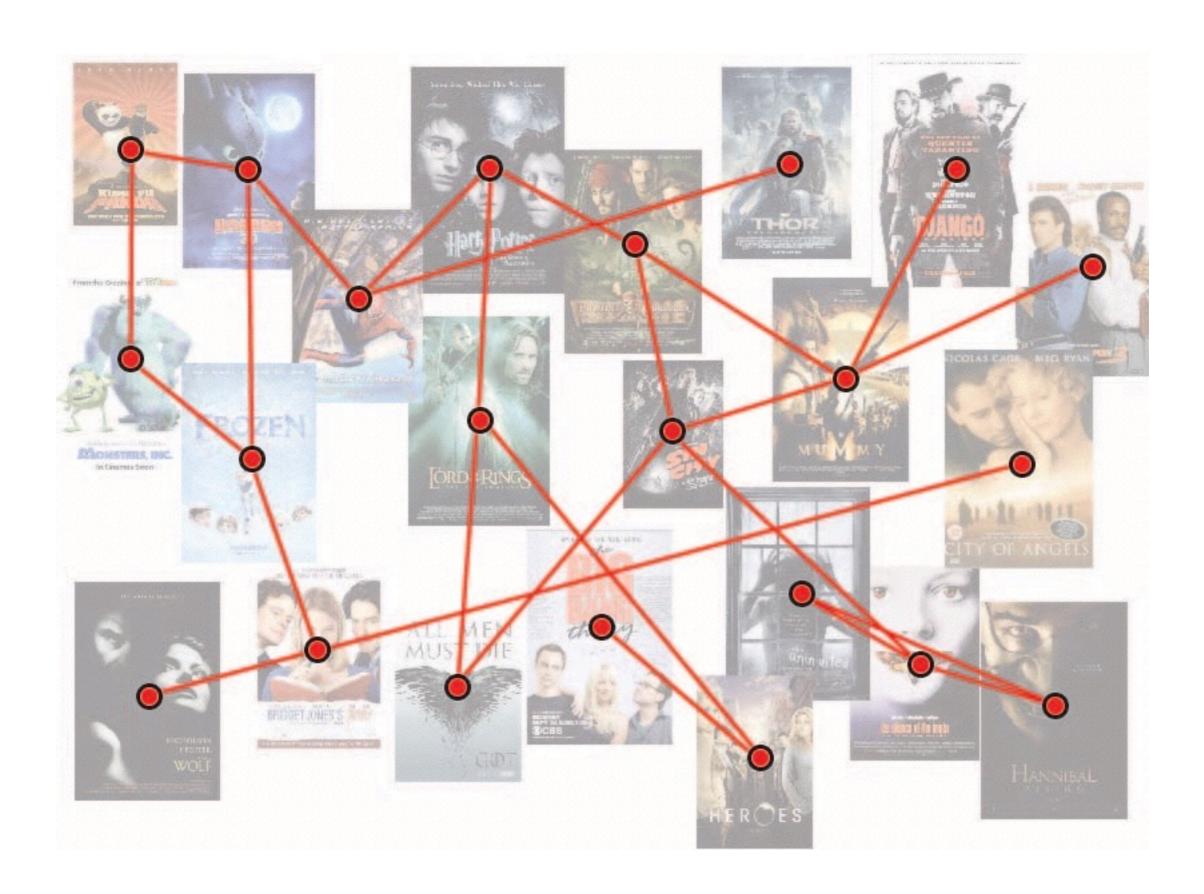


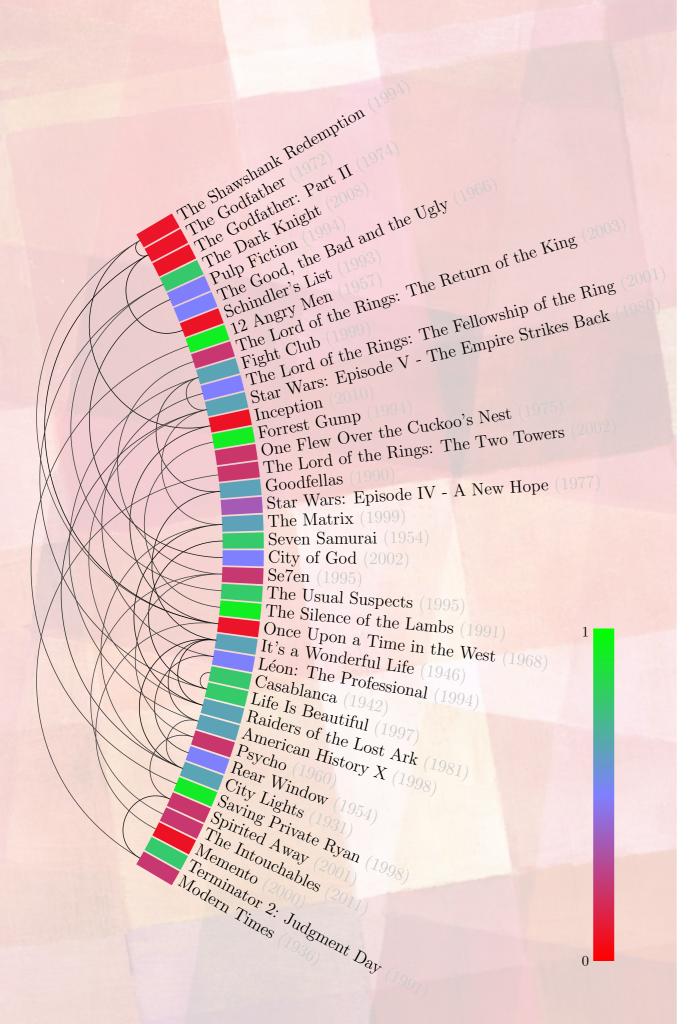
- ▶ 28. 11. 2016 by Daniele.Calandriello@inria.fr
- Content (this time lecture in class + coding at home)
 - Large-scale graph construction and processing (in class)
 - Scalable algorithms:
 - Online face recognizer (to code in Matlab)
 - Iterative label propagation (to code in Matlab)
 - Graph sparsification (presented in class)
- AR: record a video with faces
- Short written report
- Questions to piazza
- ► *Deadline:* 12. 12. 2016 (Today!)
- http://researchers.lille.inria.fr/~calandri/teaching.html

FINAL CLASS PROJECTS



- time and formatting description on the class website
- grade: report + short presentation of the team
- deadlines
 - ▶ 5. 1. 2017 final report (for all projects)
 - ▶ 9. 1. 2017, presentation in class (Cournot C102)
 - alternatively Jan 2017, remote presentations (other projects)
- project report: 5-10 pages in NIPS format
- presentation: 15+5 minutes (time it!)
- everybody has to present
- book presentation time slot on the website
- explicitly state the contributions





Example of a graph bandit problem

movie recommendation

- recommend movies to a **single** user
- goal: maximise the sum of the ratings (minimise regret)
- good prediction after just a few steps

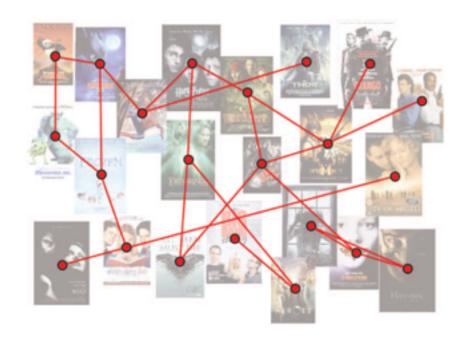
$$T \ll N$$

- extra information
 - ratings are **smooth** on a graph
- main question: can we learn **faster**?

GETTING REAL



Let's be lazy and ignore the structure



 \rightarrow

FROZEN

ALL MIN

ALL MIN

FORMANIA

#actions

#rounds

Multi-armed bandit problem!

Worst case regret (to the best fixed strategy)

 $R_T = \mathcal{O}\left(\sqrt{NT}\right)$

Matching lower bound (Auer, Cesa-Bianchi, Freund, Schapire 2002)

How big is N? Number of movies on http://www.imdb.com/stats: 4,003,294

Problem: Too many actions!

LEARNING FASTER

#actions

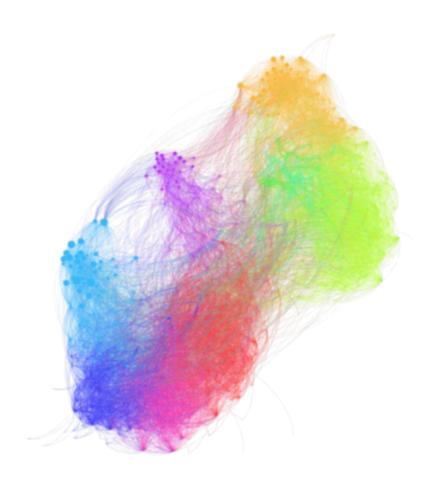


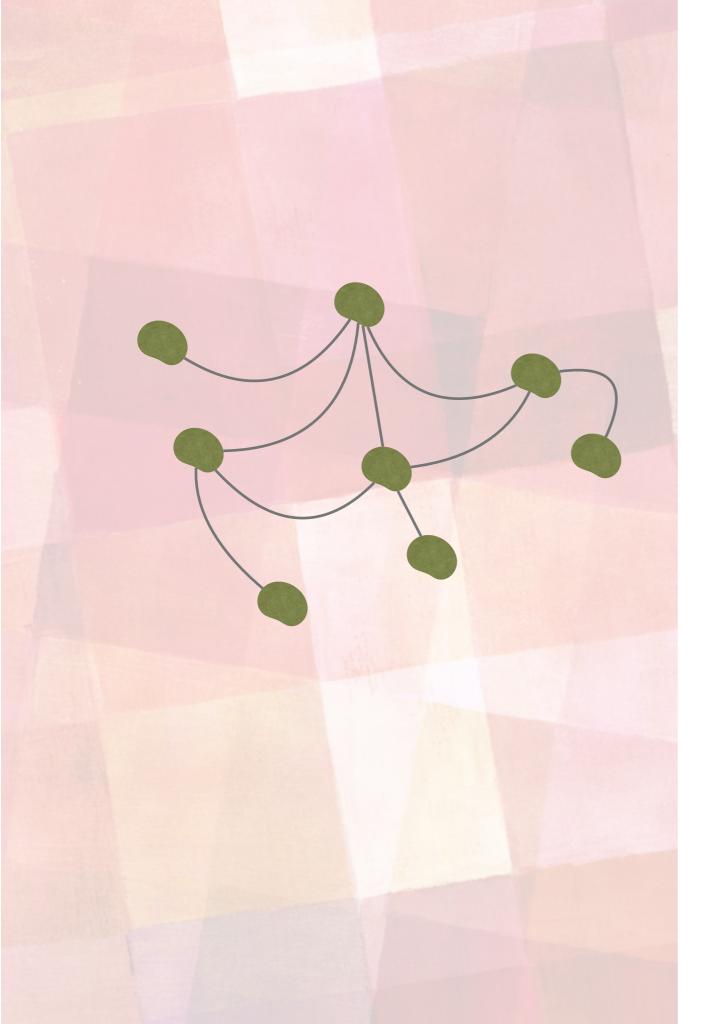
$$R_T = \mathcal{O}\left(\sqrt{NT}\right)^{\text{\#rounds}}$$

- Arm independence is too strong and unnecessary
- Replace N with something much smaller
 - problem/instance/data dependent
 - example: linear bandits N to D

#dimensions

- Today: Graph Bandits!
 - sequential problems where actions are nodes on a graph
 - find strategies that replace N with a smaller graph-dependent quantity





GRAPH BANDITS: GENERAL SETUP

Every round **t** the learner

- \triangleright picks a node $I_t \in [N]$
- \triangleright incurs a loss ℓ_{t,I_t}
- optional feedback

The performance is total expected regret

$$R_T = \max_{i \in [N]} \mathbb{E} \left[\sum_{t=1}^T (\ell_{t,I_t} - \ell_{t,i}) \right]$$

loss

Specific problems differ in 2. feedback

guarantees

STRUCTURES IN BANDIT PROBLEMS



GRAPHS

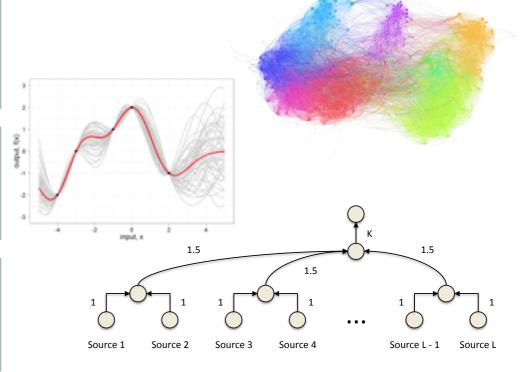
KERNELS

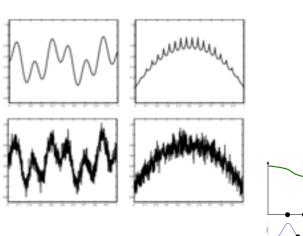
POLYMATROIDS

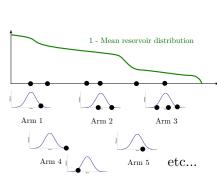
BLACK-BOX FUNCTIONS

STRUCTURES WITHOUT TOPOLOGY

. . .







SPECIFIC GRAPH BANDIT SETTINGS



smoothness spectral bandits

$$R_T = \widetilde{\mathcal{O}}\left(\frac{d}{\sqrt{T \ln T}}\right)$$

independence number

side observations on graphs

$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha} T \ln N}\right)$$

#relevant eigenvectors

> influence maximisation revealing bandits

$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{r_* T D_*}\right)$$

detectable dimension

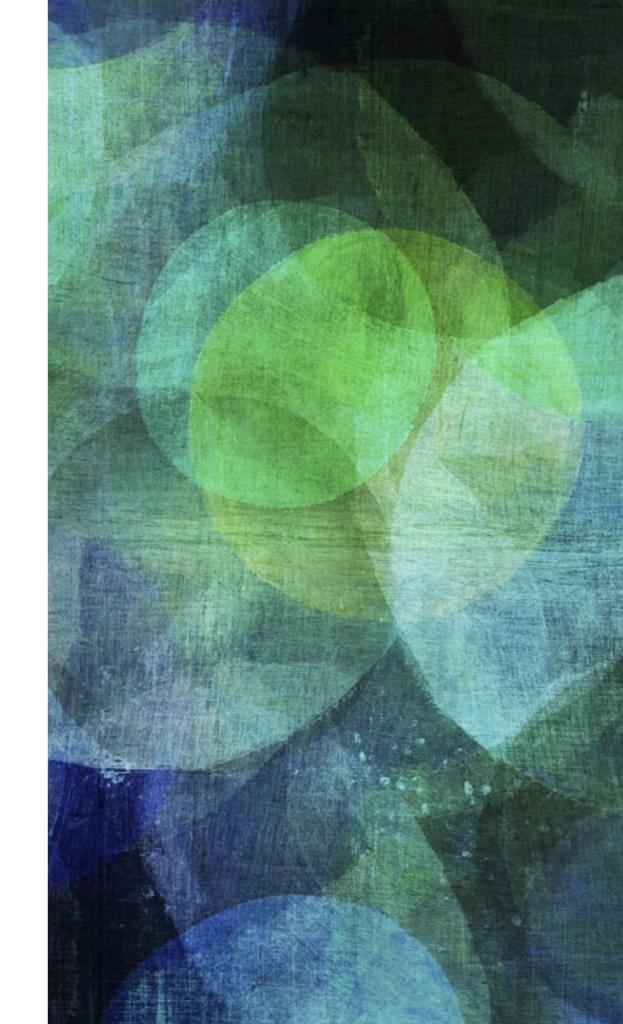
noisy side observations

on graphs
$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{\alpha^* T \ln N}\right)$$

effective independence number MV, Munos, Kveton, Kocák: **Spectral Bandits for Smooth Graph Functions**, ICML 2014 Kocák, MV, Munos, Agrawal: **Spectral Thompson Sampling**, AAAI 2014 Hanawal, Saligrama, MV, Munos: **Cheap Bandits**, ICML 2015

SPECTRAL BANDITS

exploiting smoothness of rewards on graphs



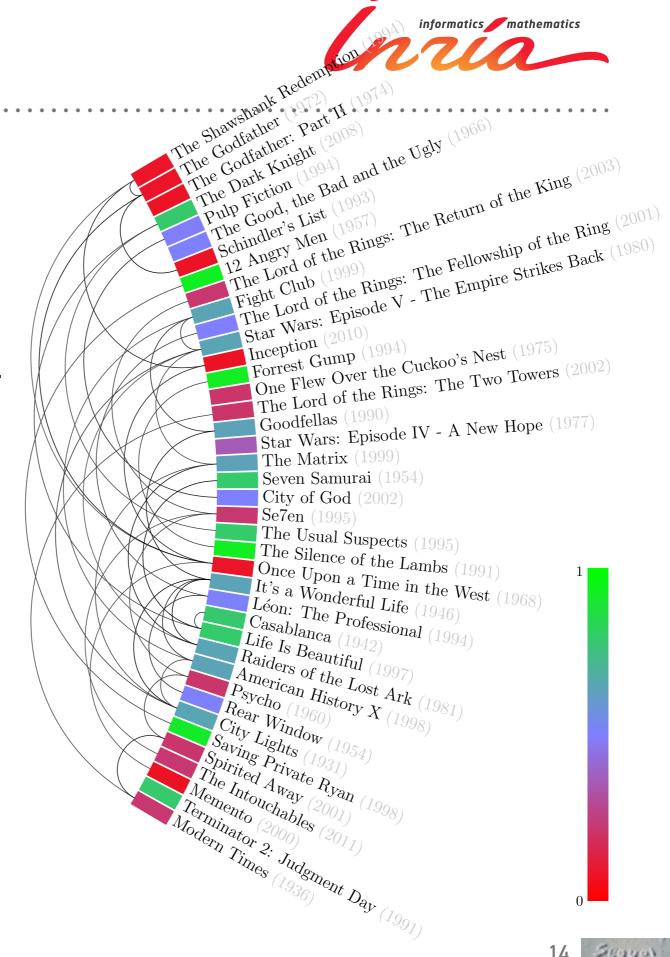
SPECTRAL BANDITS

Assumptions

- ▶ Unknown reward function $f:V(G)\to \mathbb{R}$.
- Function f is **smooth** on a graph.
- Neighboring movies \Rightarrow similar preferences.
- \triangleright Similar preferences \Rightarrow neighboring movies.

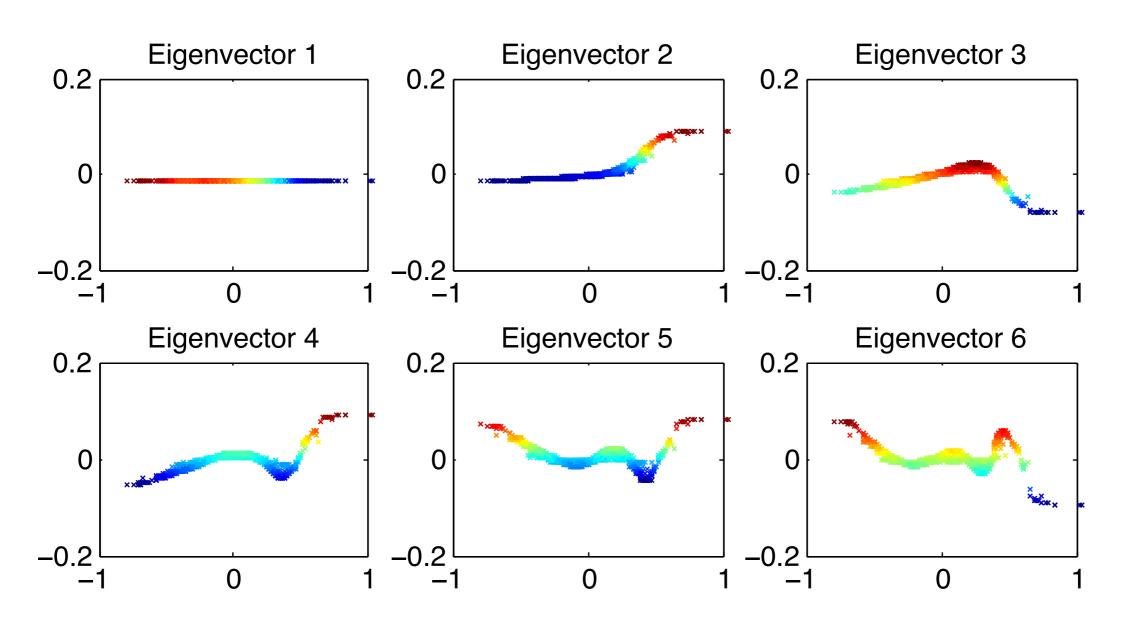
Desiderata

An algorithm useful in the case $T \ll N!$



FLIXSTER DATA





Eigenvectors from the Flixster data corresponding to the smallest few eigenvalues of the graph Laplacian projected onto the first principal component of data. Colors indicate the values.

SPECTRAL BANDIT: LEARNING SETTING



Learning setting for a bandit algorithm π

- ▶ In each time t step choose a node $\pi(t)$.
- ▶ the $\pi(t)$ -th row $\mathbf{x}_{\pi(t)}$ of the matrix \mathbf{Q} corresponds to the arm $\pi(t)$.
- ▶ Obtain noisy reward $r_t = \mathbf{x}_{\pi(t)}^\mathsf{T} \alpha^* + \varepsilon_t$. Note: $\mathbf{x}_{\pi(t)}^\mathsf{T} \alpha^* = f_{\pi(t)}$
 - ▶ ε_t is R-sub-Gaussian noise. $\forall \xi \in \mathbb{R}, \mathbb{E}[e^{\xi \varepsilon_t}] \leq \exp(\xi^2 R^2/2)$
- Minimize cumulative regret

$$R_T = T \max_{a} (\mathbf{x}_a^{\mathsf{T}} \boldsymbol{\alpha}^*) - \sum_{t=1}^{T} \mathbf{x}_{\pi(t)}^{\mathsf{T}} \boldsymbol{\alpha}^*.$$

Can we just use linear bandits?

LINEAR VS. SPECTRAL BANDITS



Linear bandit algorithms

- ► LinUCB
 - Regret bound $\approx D\sqrt{T \ln T}$
- LinearTS
 - Regret bound $\approx D\sqrt{T \ln N}$

(Li et al., 2010)

(Agrawal and Goyal, 2013)

Note: D is ambient dimension, in our case N, length of x_i . Number of actions, e.g., all possible movies \rightarrow **HUGE!**

- Spectral bandit algorithms
 - SpectralUCB
 - Regret bound $\approx d\sqrt{T \ln T}$
 - \triangleright Operations per step: D^2N
 - SpectralTS
 - Regret bound $\approx d\sqrt{T \ln N}$
 - ▶ Operations per step: $D^2 + DN$

Note: d is effective dimension, usually much smaller than D.

(Valko et al., ICML 2014)

(Kocák et al., AAAI 2014)

SPECTRAL BANDITS - EFFECTIVE DIMENSION



▶ **Effective dimension:** Largest *d* such that

$$(d-1)\lambda_d \leq \frac{T}{\log(1+T/\lambda)}.$$

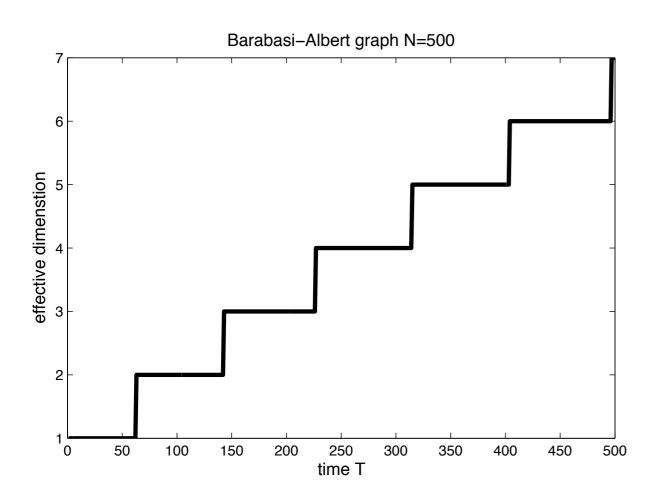
- Function of time horizon and graph properties
- λ_i : *i*-th smallest eigenvalue of **Λ**.
- \triangleright λ : Regularization parameter of the algorithm.

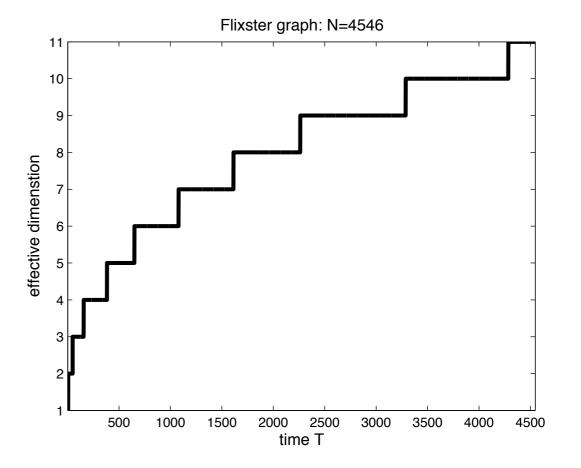
Properties:

- \triangleright d is small when the coefficients λ_i grow rapidly above time.
- d is related to the number of "non-negligible" dimensions.
- Usually d is much smaller than D in real world graphs.
- Can be computed beforehand.

SPECTRAL BANDITS - EFFECTIVE DIMENSION





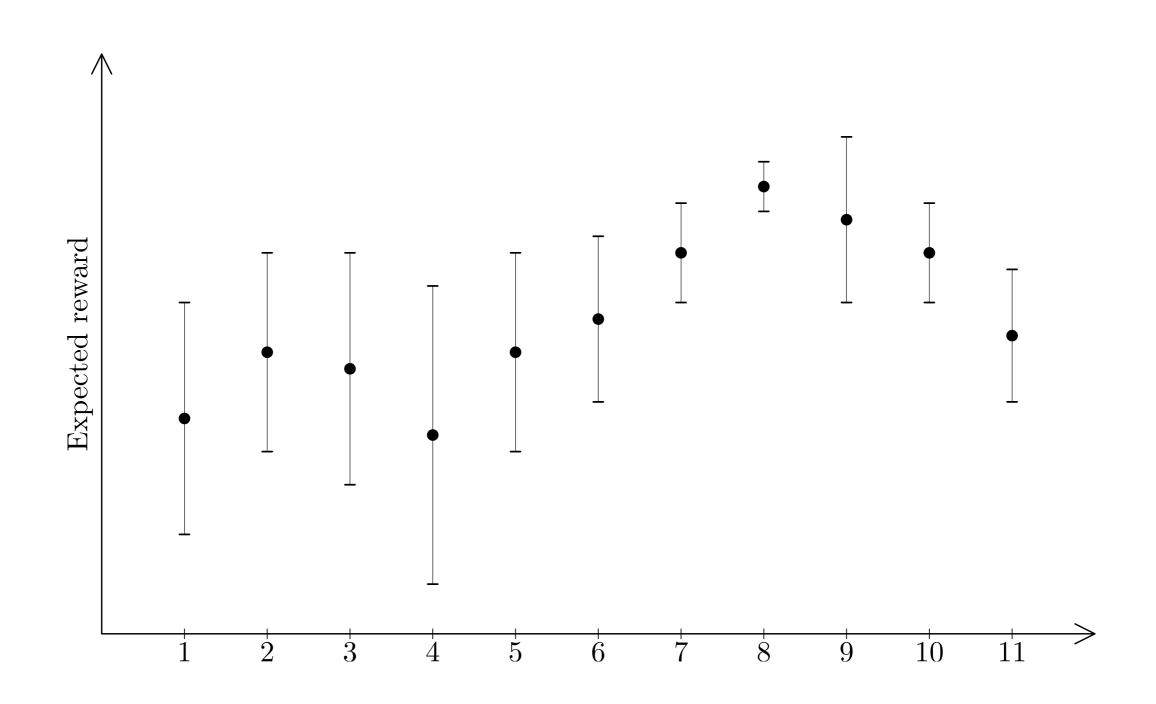


$$d \ll D$$

Note: In our setting T < N = D.

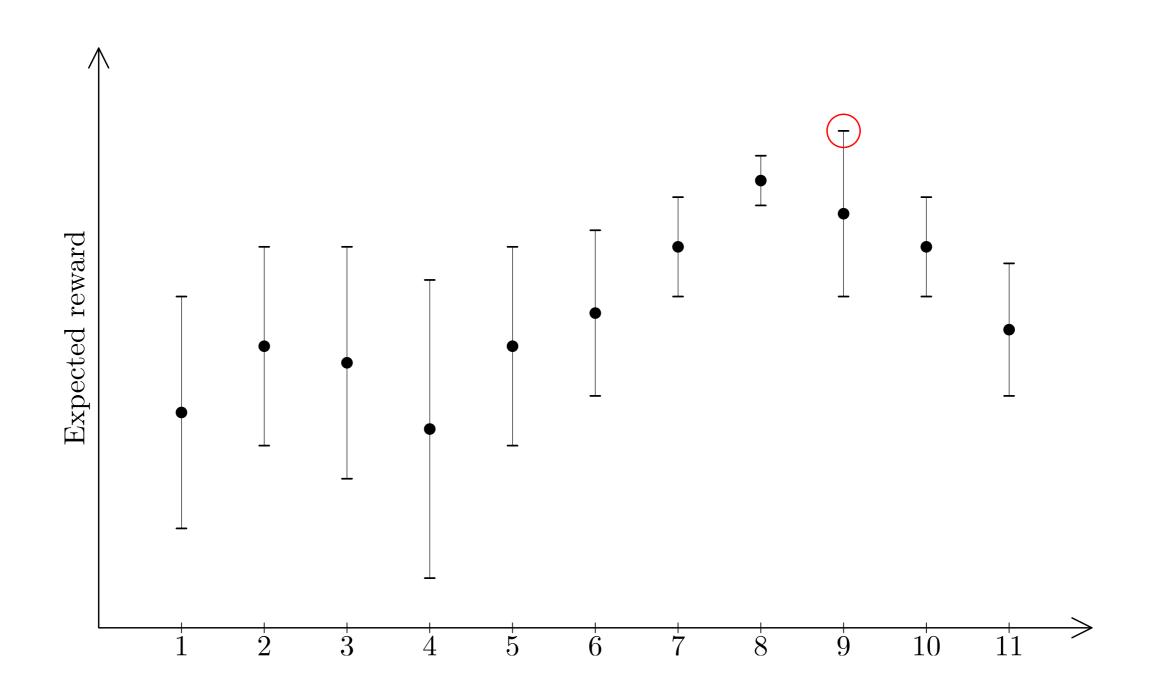
UPPER CONFIDENCE BOUND BASED ALGOS





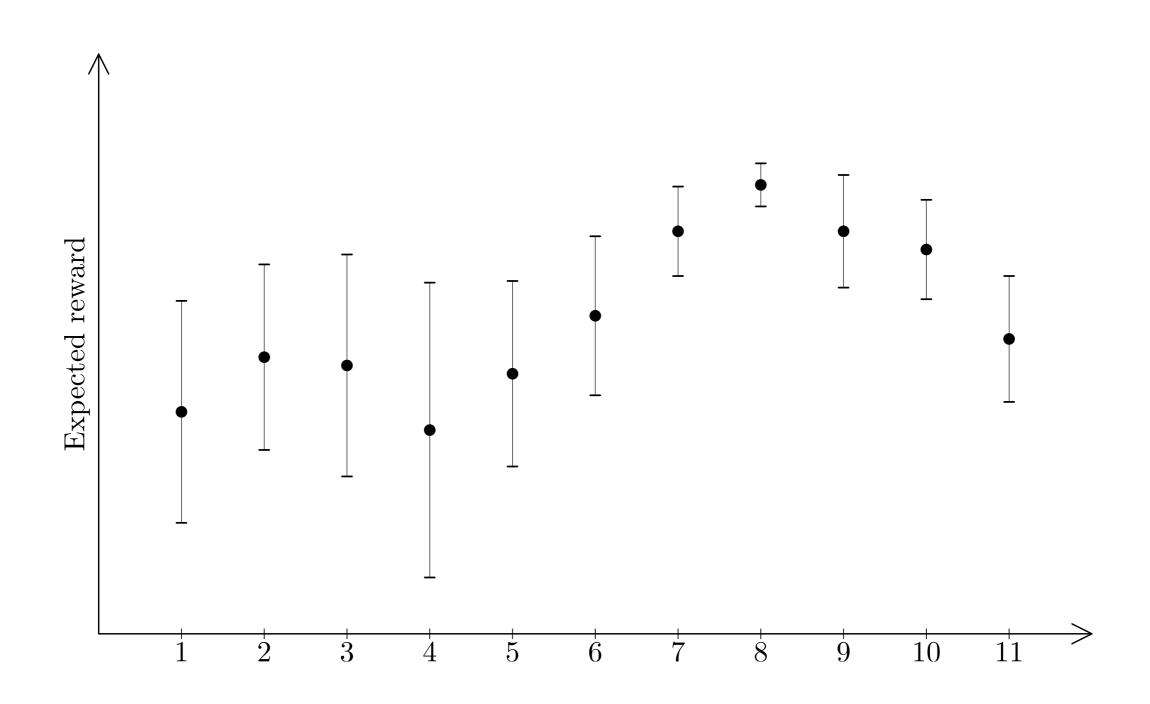
UPPER CONFIDENCE BOUND BASED ALGOS





UPPER CONFIDENCE BOUND BASED ALGOS





SPECTRAL UCB



Given a vector of weights α , we define its Λ norm as

$$\|\alpha\|_{\mathbf{\Lambda}} = \sqrt{\sum_{k=1}^{N} \lambda_k \alpha_k^2} = \sqrt{\alpha^{\mathsf{T}} \mathbf{\Lambda} \alpha},$$

and fit the ratings r_v with a (regularized) least-squares estimate

$$\widehat{\alpha}_t = \operatorname*{arg\,min}_{\pmb{lpha}} \left(\sum_{v=1}^t \left[\langle \mathbf{x}_v, \pmb{lpha}
angle - r_v
ight]^2 + \| \pmb{lpha} \|_{\pmb{\Lambda}}^2
ight).$$

 $\|\alpha\|_{\Lambda}$ is a penalty for non-smooth combinations of eigenvectors.

SPECTRAL UCB

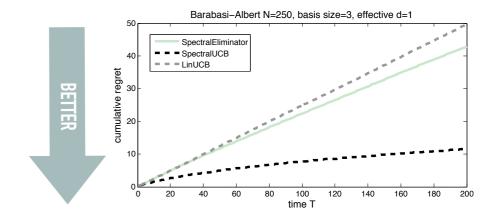


```
1: Input:
  2: N, T, \{\Lambda_L, \mathbf{Q}\}, \lambda, \delta, R, C
  3: Run:
         \Lambda \leftarrow \Lambda_1 + \lambda I
         d \leftarrow \max\{d: (d-1)\lambda_d \leq T/\ln(1+T/\lambda)\}
  6: for t = 1 to T do
              Update the basis coefficients \widehat{\alpha}:
         \mathbf{X}_t \leftarrow [\mathbf{x}_{\pi(1)}, \dots, \mathbf{x}_{\pi(t-1)}]^{\mathsf{T}}
  8:
        \mathbf{r} \leftarrow [r_1, \dots, r_{t-1}]^{\mathsf{T}}
10: \mathbf{V}_t \leftarrow \mathbf{X}_t \mathbf{X}_t^\mathsf{T} + \mathbf{\Lambda}
11: \widehat{\boldsymbol{\alpha}}_t \leftarrow \mathbf{V}_t^{-1} \mathbf{X}_t^{\mathsf{T}} \mathbf{r}
12: c_t \leftarrow 2R\sqrt{d\ln(1+t/\lambda)+2\ln(1/\delta)}+C
          \pi(t) \leftarrow \operatorname{arg\,max}_{a} \left( \mathbf{x}_{a}^{\mathsf{T}} \widehat{\alpha} + c_{t} \| \mathbf{x}_{a} \|_{\mathbf{V}_{\star}^{-1}} \right)
13:
14:
              Observe the reward r_t
15: end for
```

SPECTRALUCB REGRET BOUND

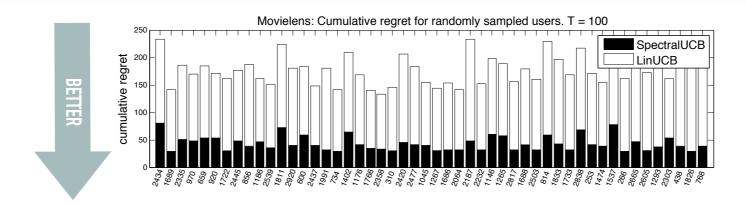


- ▶ *d*: Effective dimension.
- \triangleright λ: Minimal eigenvalue of $\Lambda = \Lambda_L + \lambda I$.
- ightharpoonup C: Smoothness upper bound, $\|\alpha^*\|_{\Lambda} \leq C$.
- $ightharpoonup \mathbf{x}_i^{\mathsf{T}} oldsymbol{lpha}^* \in [-1,1]$ for all i.



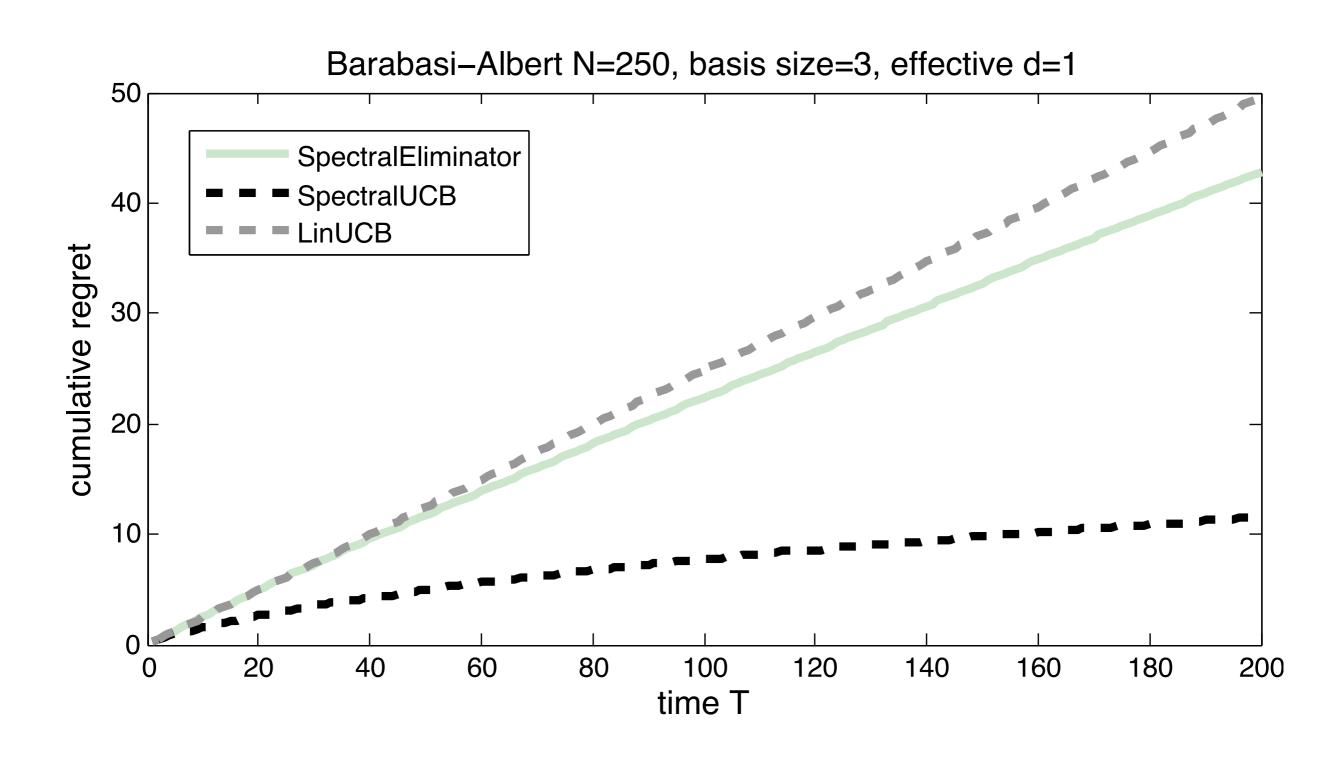
The **cumulative regret** R_T of **SpectralUCB** is with probability $1 - \delta$ bounded as

$$R_T \le \left(8R\sqrt{d\ln\frac{\lambda+T}{\lambda}+2\ln\frac{1}{\delta}}+4C+4\right)\sqrt{dT\ln\frac{\lambda+T}{\lambda}}.$$



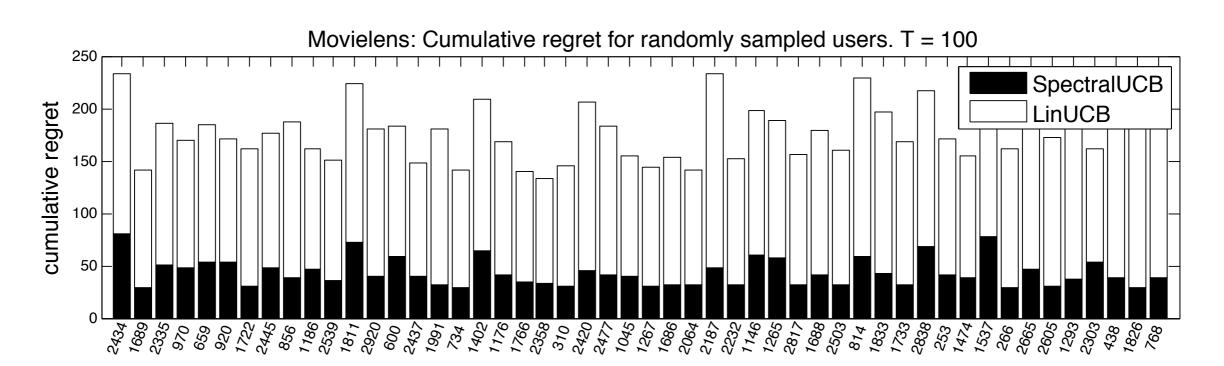
SPECTRAL UCN ON BA GRAPH

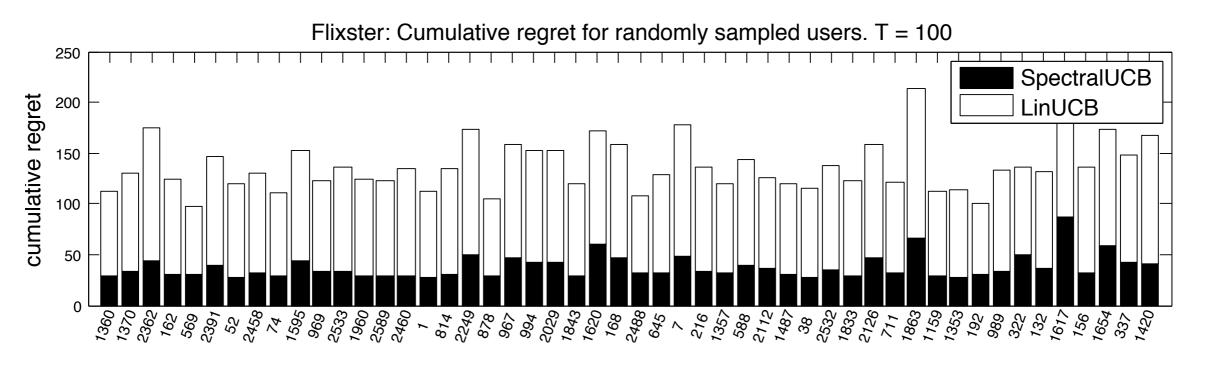




SPECTRAL UCN ON REAL DATA





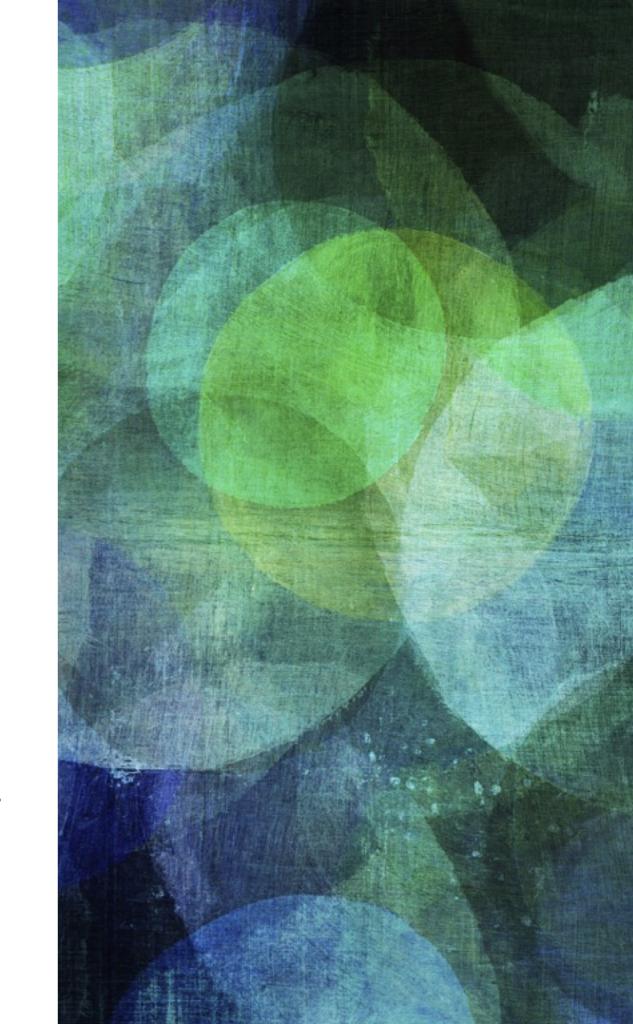


Kocák, Neu, MV, Munos: **Efficient learning by implicit exploration in bandit problems** with side observations, NIPS 2014

Kocák, Neu, MV: **Online learning with Erdos-Rényi side-observation graphs** UAI 2016 (to appear)

GRAPH BANDITS WITH SIDE OBSERVATIONS

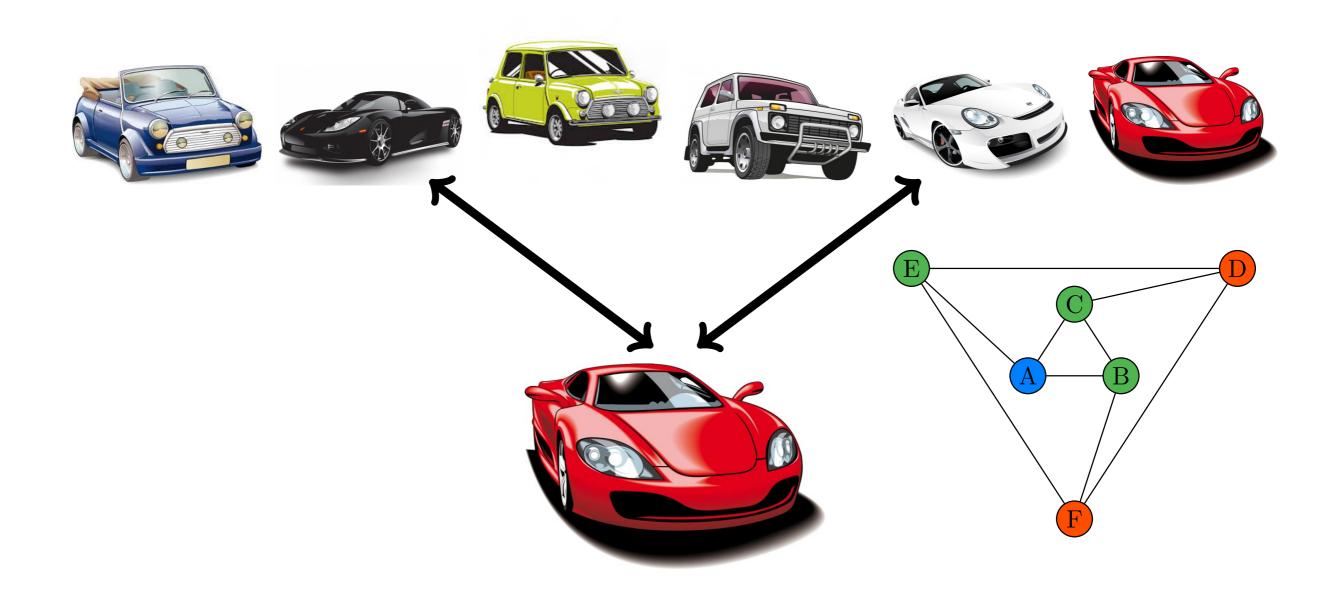
exploiting free observations from neighbouring nodes



SIDE OBSERVATIONS: UNDIRECTED



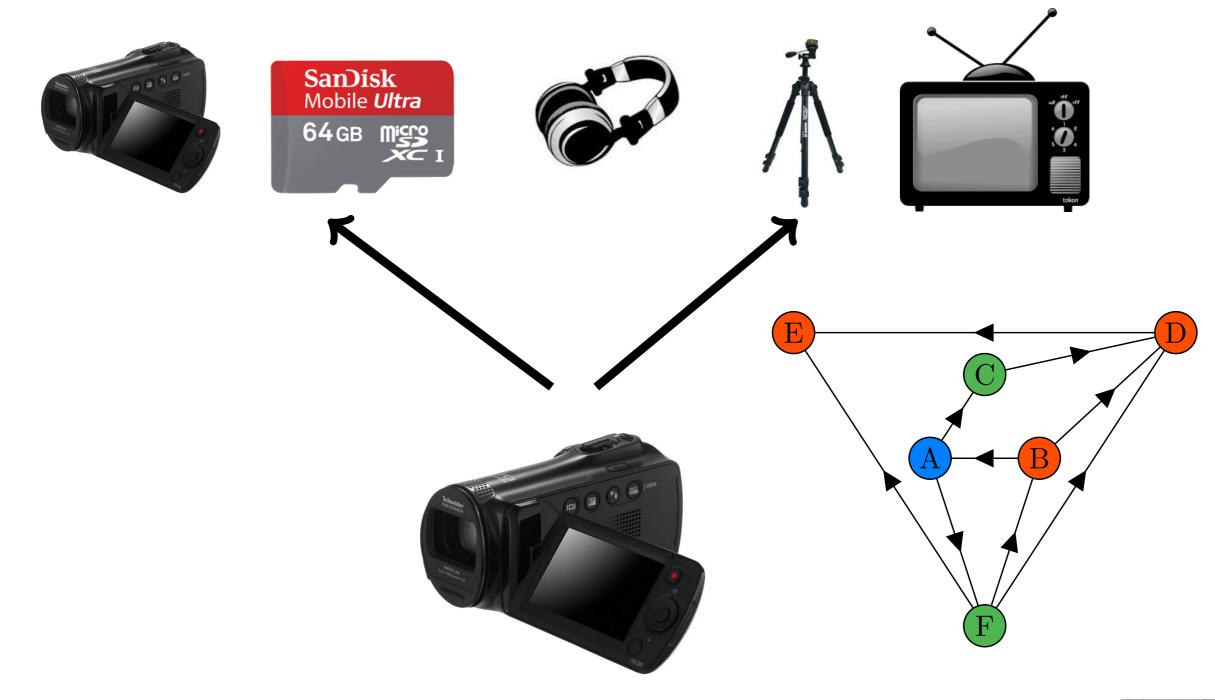
Example 1: undirected observations



SIDE OBSERVATIONS: DIRECTED



Example 2: Directed observation



SIDE OBSERVATIONS - AN INTERMEDIATE GAME COLOR



Full-information

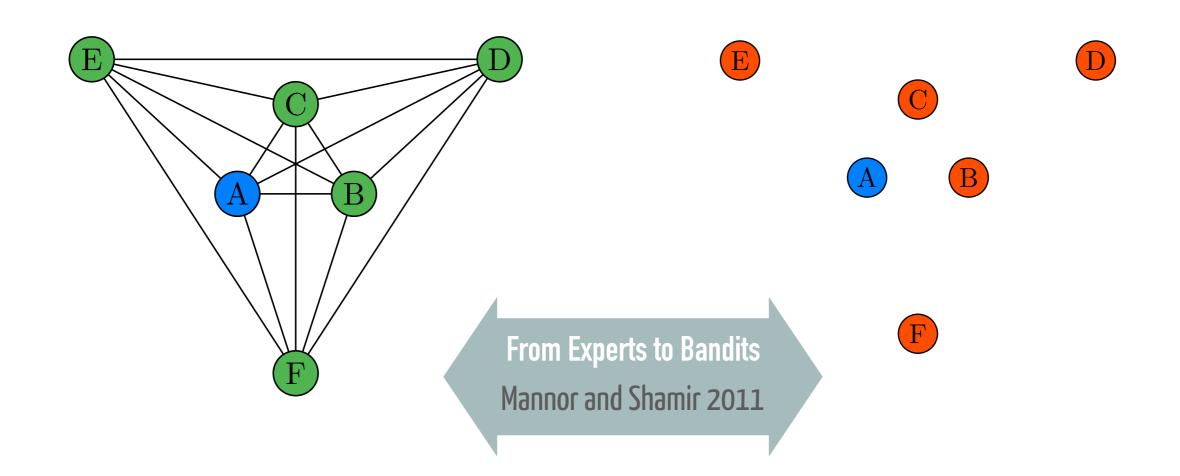
- observe losses of all actions
- example: Hedge

$$R_T = \widetilde{\mathcal{O}}(\sqrt{T})$$

Bandits

- observe losses of the chosen action
- example: EXP3

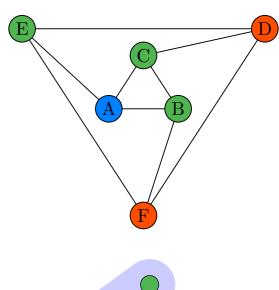
$$R_T = \widetilde{\mathcal{O}}(\sqrt{NT})$$

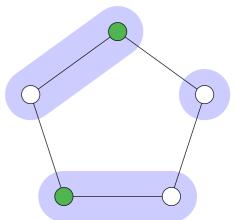


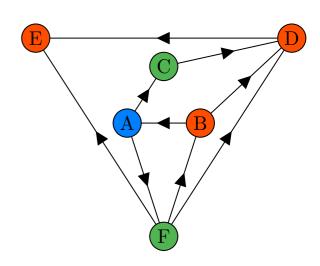
KNOWLEDGE OF OBSERVATION GRAPHS



- ELP (Mannor and Shamir 2011)
 - EXP3 with "LP balanced exploration"
 - undirected $O(\sqrt{(\alpha T)}) = -\text{needs to know } G_t$
 - directed case $O(\sqrt{(cT)})$ needs to know G_t
- EXP3-SET (Alon, Cesa-Bianchi, Gentile, Mansour, 2013)
 - undirected $O(\sqrt{(\alpha T)})$ \square does not need to know G_t \square
- EXP3-DOM (Alon, Cesa-Bianchi, Gentile, Mansour, 2013)
 - directed $O(\sqrt{(\alpha T)}) \square \text{need to know } G_t$
 - calculates dominating set







EXP3-IX: IMPLICIT EXPLORATION



Algorithm 1 EXP3-IX

- 1: **Input:** Set of actions S = [d],
- 2: parameters $\gamma_t \in (0,1), \eta_t > 0$ for $t \in [T]$.
- 3: **for** t = 1 **to** T **do**
- 4: $w_{t,i} \leftarrow (1/d) \exp(-\eta_t \widehat{L}_{t-1,i}) \text{ for } i \in [d]$
- 5: An adversary privately chooses losses $\ell_{t,i}$ for $i \in [d]$ and generates a graph G_t
- 6: $W_t \leftarrow \sum_{i=1}^d w_{t,i}$
- 7: $p_{t,i} \leftarrow w_{t,i}/W_t$
- 8: Choose $I_t \sim p_t = (p_{t,1}, \dots, p_{t,d})$
- 9: Observe graph G_t
- 10: Observe pairs $\{i, \ell_{t,i}\}$ for $(I_t \to i) \in G_t$
- 11: $o_{t,i} \leftarrow \sum_{(j \to i) \in G_t} p_{t,j} \text{ for } i \in [d]$
- 12: $\hat{\ell}_{t,i} \leftarrow \frac{\ell_{t,i}}{o_{t,i} + \gamma_t} \mathbb{1}_{\{(I_t \to i) \in G_t\}} \text{ for } i \in [d]$
- 13: **end for**

Benefits of the implicit exploration

- no need to know the graph before
- no need to estimate dominating set
- no need for doubling trick
- no need for aggregation

$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha}\,T\,\ln N}\right)$$

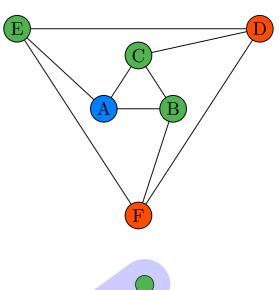
Optimistic bias for the loss estimates

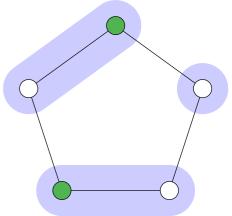
$$\mathbb{E}[\hat{\ell}_{t,i}] = \frac{\ell_{t,i}}{o_{t,i} + \gamma} o_{t,i} + 0(1 - o_{t,i}) = \ell_{t,i} - \ell_{t,i} \frac{\gamma}{o_{t,i} + \gamma} \leq \ell_{t,i}$$

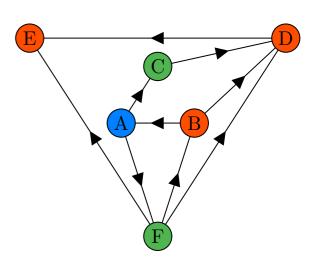
FOLLOW UPS



- EXP3-IX (Kocák, Neu, MV, Munos, 2014)
 - directed $O(\sqrt{(\alpha T)})$ does not need to know G_t
- EXP3.G (Alon, Cesa-Bianchi, Dekel, Koren, 2015)
 - directed $O(\sqrt{(\alpha T)})$ does not need to know G_t
 - mixes uniform distribution
 - more general algorithm for settings **beyond bandits**
 - high-probability bound
- Neu 2015: high-probability bound for EXP3-IX



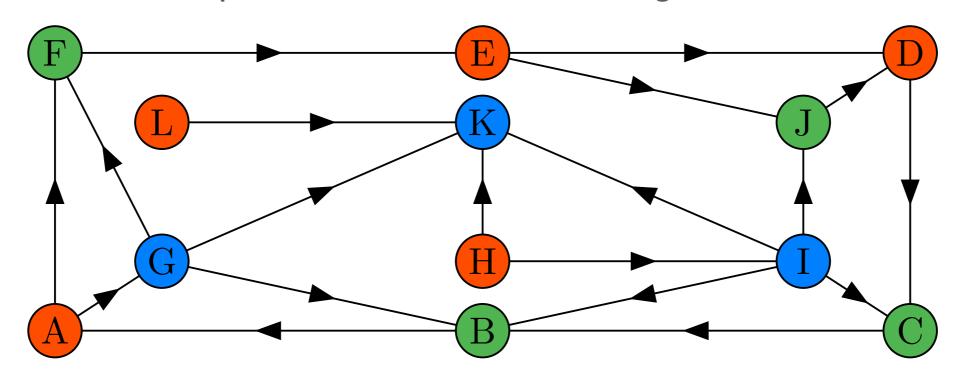




COMPLEX GRAPH ACTIONS



Example: online shortest path semi-bandits with observing traffic on the side streets

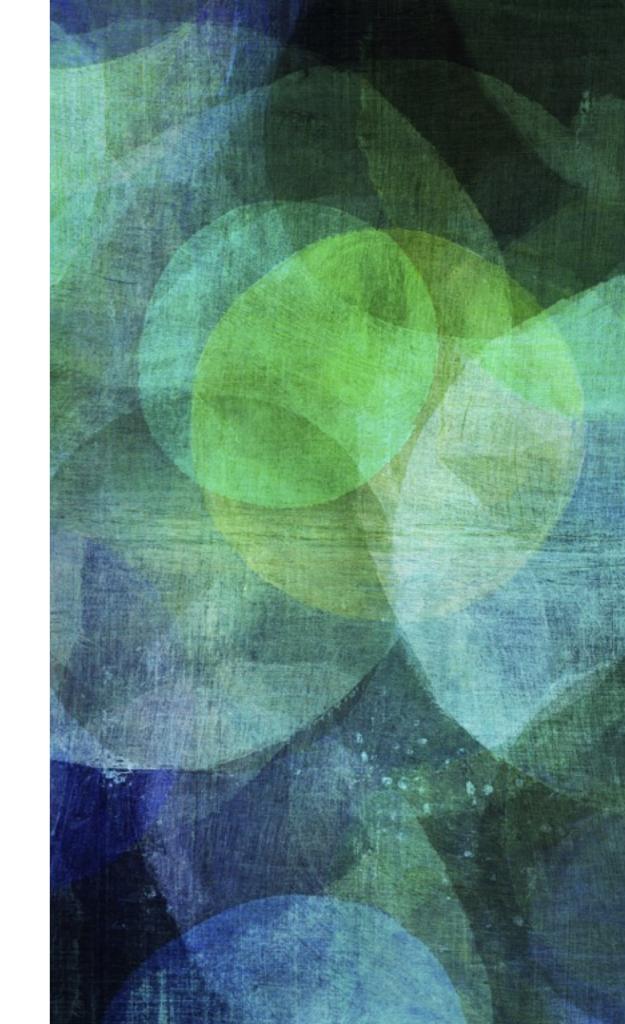


- ▶ Play action $V_t \in S \subset \{0,1\}^N$, $\|\mathbf{v}\|_1 \leq m$ from all $\mathbf{v} \in S$
- ightharpoonup Obtain losses $\mathbf{V}_t^{\mathsf{T}} \ell_t$
- Observe additional losses according to the graph

$$R_T = \widetilde{\mathcal{O}}\left(m^{3/2}\sqrt{\sum_{t=1}^T \alpha_t}\right) = \widetilde{\mathcal{O}}\left(m^{3/2}\sqrt{\overline{\alpha}T}\right)$$

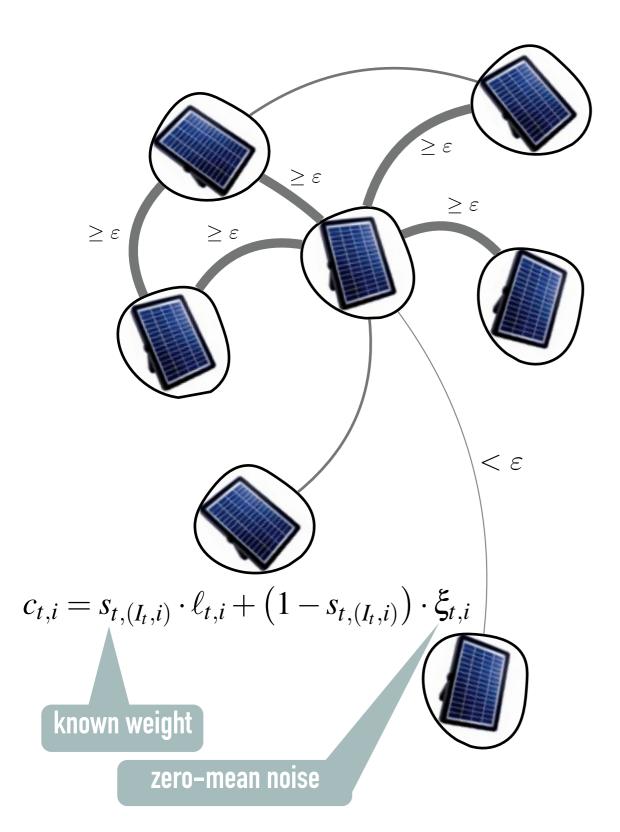
GRAPH BANDITS WITH NOISY SIDE OBSERVATIONS

exploiting side observations that can be perturbed by certain level of noise



NOISY SIDE OBSERVATIONS





Want: only reliable information!

- 1) If we know the perfect cutoff **E**
- reliable: use as exact
- unreliable: rubbish then we can improve over pure bandit setting!
- 2) Treating noisy observation induces bias

What can we hope for?

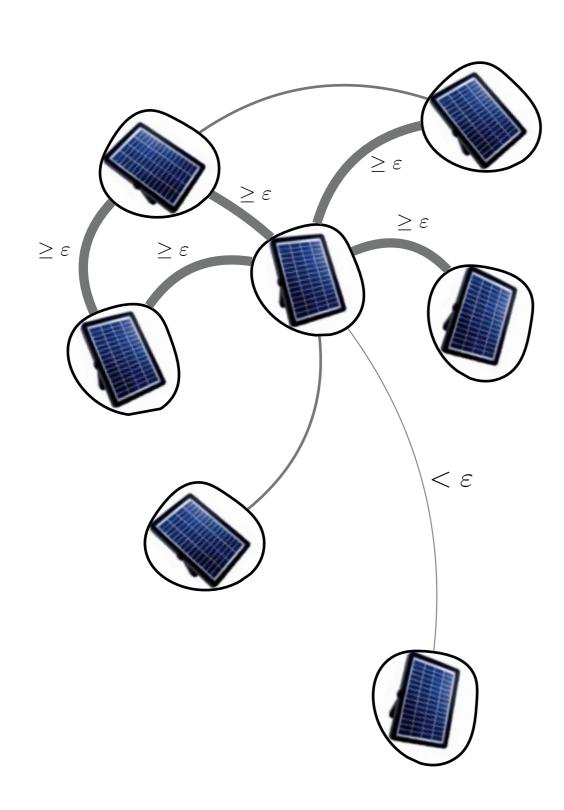
$$\widetilde{\mathcal{O}}\left(\sqrt{1T}\right) \leq \leq \widetilde{\mathcal{O}}\left(\sqrt{NT}\right)$$

effective independence number

Can we learn without knowing either ε or α^* ?

NOISY SIDE OBSERVATIONS





Threshold estimate
$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{\overline{lpha}^{\star}T}
ight)$$
 $\widehat{\ell}^{(\mathrm{T})} = \underbrace{c_{t,i}\mathbb{I}_{\left\{s_{t,(I_t,i)}\geq arepsilon_t
ight\}}}$

$$\widehat{\ell}_{t,i}^{(\mathrm{T})} = \frac{c_{t,i} \mathbb{I}_{\left\{s_{t,(I_t,i)} \geq \varepsilon_t\right\}}}{\sum_{j=1}^{N} p_{t,j} s_{t,(j,i)} \mathbb{I}_{\left\{s_{t,(j,i)} \geq \varepsilon_t\right\}} + \gamma_t}$$

WIX estimate

$$R_T = \widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha}^*T}\right)$$

$$\widehat{\ell}_{t,i} = \frac{s_{t,(I_t,i)} \cdot c_{t,i}}{\sum_{j=1}^{N} p_{t,j} s_{t,(j,i)}^2 + \gamma_t}$$

Since $\alpha^* \leq \alpha(1)/1 \leq N$

incorporating noisy observations does not hurt

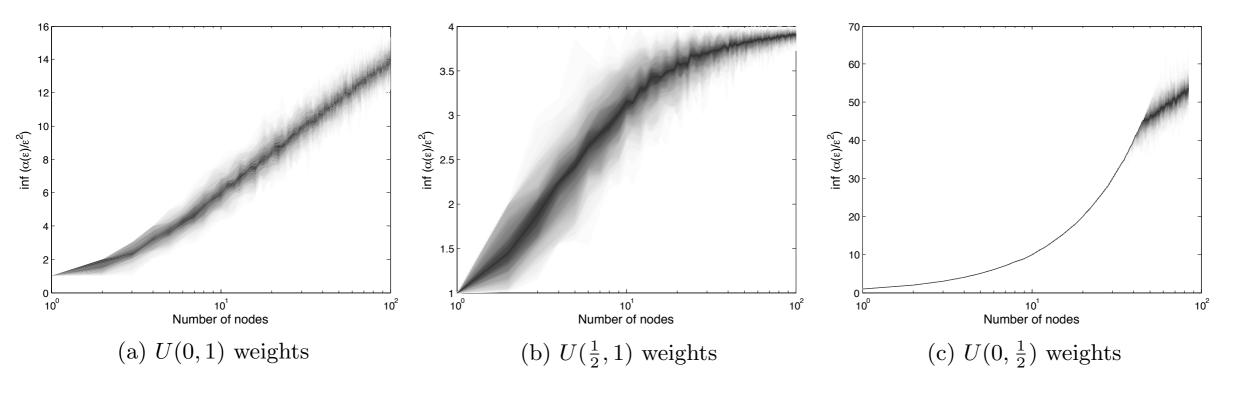
$$\widetilde{\mathcal{O}}\left(\sqrt{\overline{\alpha}^{\star}T}\right) \leq \widetilde{\mathcal{O}}\left(\sqrt{NT}\right)$$

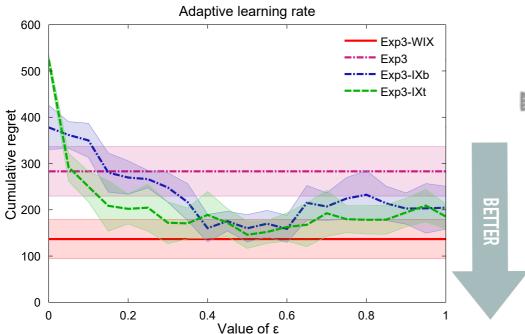
But how much does it help?

EMPIRICAL RESULTS



EMPIRICAL α^* FOR RANDOM GRAPHS WITH IID WEIGHTS





- **special case:** if s_{ij} is either 0 or ϵ than $\alpha *= \alpha/\epsilon^2$
 - For this special case, there is a matches $\Theta(\sqrt{(\alpha T)/\epsilon})$ by Wu, György, Szepesvári, 2015.



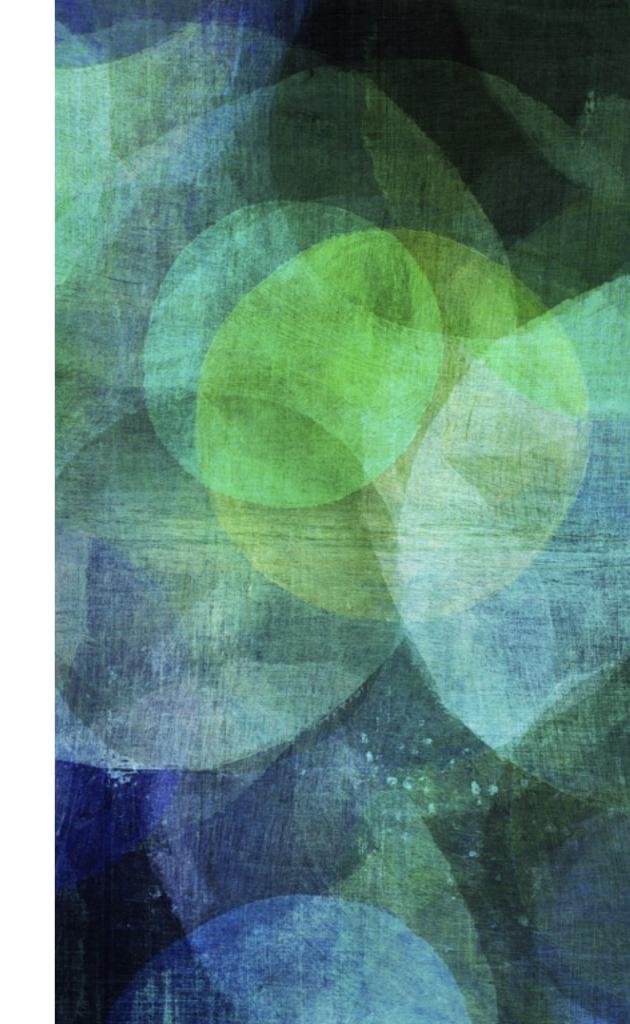
NEW DIRECTIONS: UNKNOWN GRAPHS!

- Learning on the graph while learning the graph?
 - most of algorithms require (some) knowledge of the graph
 - not always available to the learner
- Question: Can we learn faster without knowing the graphs?
 - example: social network provider has little incentive to reveal the graphs to advertisers
- Answer: **Cohen, Hazan, and Koren:** Online learning with **feedback** graphs without the graphs (ICML June 19-24, 2016)
 - **NO!** (in general we cannot, but possible in the stochastic case)
- Coming up next:
 - Erdös-Rényi side observation graphs (UAI June 25-26, 2016)
 - Influence Maximisation (AISTATS 2016 + https://arxiv.org/abs/1605.06593)

Kocák, Neu, MV: **Online learning with Erdos-Rényi side-observation graphs** UAI 2016 (to appear)

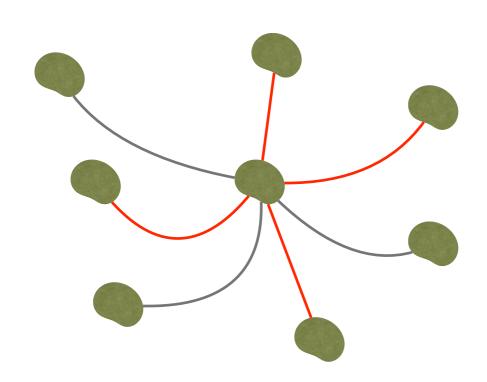
GRAPH BANDITS WITH ERDÖS-RÉNYI OBSERVATIONS

side observations from graph generators



PROTOCOL FOR ERDÖS-RÉNYI GRAPHS





is loss of i observed?

$$\widehat{\ell}_{t,i}^{\star} = \frac{O_{t,i}\ell_{t,i}}{p_{t,i} + (1 - p_{t,i})r_t}$$

probability of picking i

probability of side observation

true loss

Every round t the learner

- picks a node It
- suffers loss for It
- receives feedback
 - for It
 - for every other node with probability rt

Regret of Exp3-SET (Alon et al. 2013):

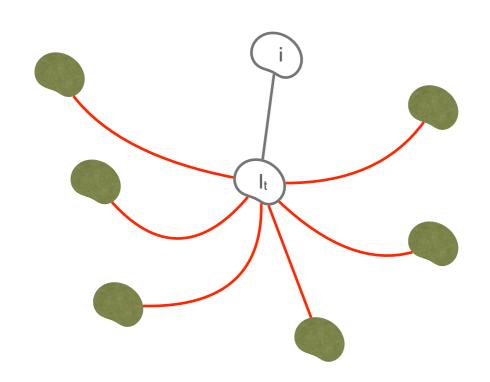
$$\mathcal{O}\left(\sqrt{\sum_t (1/r_t)(1-(1-r_t)^N)\log N}\right)$$

How to estimate **r**_t in every round when it is **changing**?

How to estimate losses without the knowledge of rt?

PROTOCOL FOR ERDÖS-RÉNYI GRAPHS





N-2 samples from Bernoulli(r_t) ... R(k)

- \triangleright N-2 samples from p_{ti} ... P(k)
- P(k) = P(k) + (1-P(k))R(k)
- $G_{ti} = min\{k : O'(k) = 1\} U \{N-1\}$
- $\mathbf{E}[\mathsf{G}_{\mathsf{ti}}] pprox \mathbf{1}/(\mathsf{p}_{\mathsf{ti}} + (\mathbf{1} \mathsf{p}_{\mathsf{ti}}) \mathsf{r}_{\mathsf{t}})$ $\widehat{\ell}_{t,i} = G_{t,i} O_{t,i} \ell_{t,i}$

is loss of i observed?

true loss

$$\widehat{\ell}_{t,i}^{\star} = \frac{O_{t,i}\ell_{t,i}}{p_{t,i} + (1 - p_{t,i})r_t}$$

probability of picking i

probability of side observation

If $r_t \ge (\log T)/(2N-2)$ then

$$\mathcal{O}\left(\sqrt{\log N \sum_{t=1}^{T} \frac{1}{r_t}}\right)$$

Lower bound (Alon et al. 2013) $\Omega(\sqrt{T/r})$

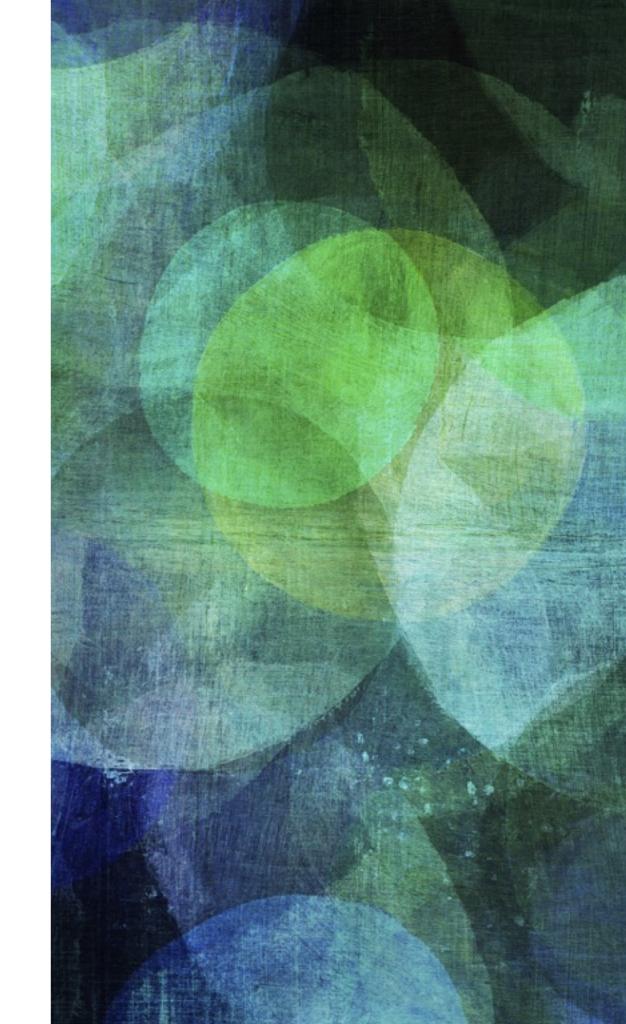
Get rid of $r_t \ge (\log T)/(2N-2)$?

Carpentier, MV: Revealing Graph Bandits for Maximising Local Influence, AISTATS 2016

Wen, Kveton, MV: Influence Maximization with Semi-Bandit Feedback, (arXiv:1605.06593)

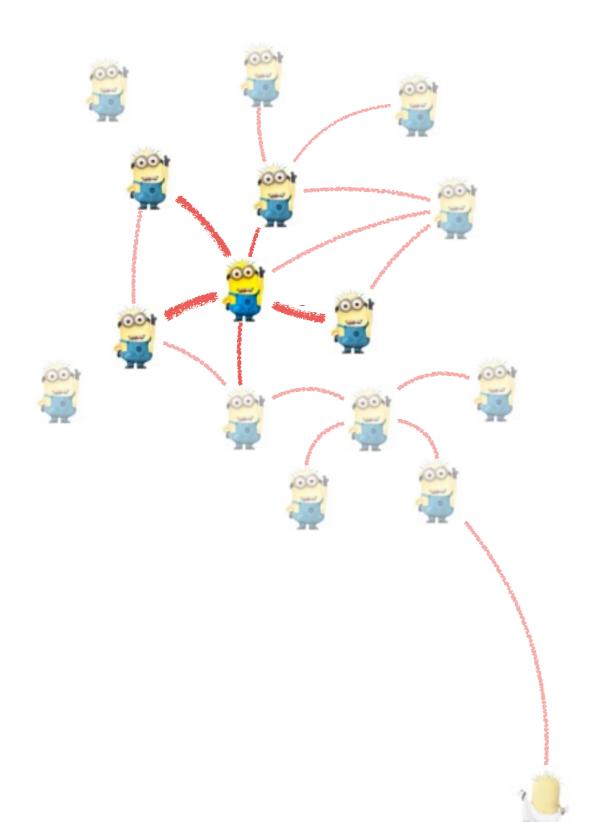
INFLUENCE MAXIMISATION

looking for the influential nodes while exploring the graph



MAXIMIZING INFLUENCE





Product placement

- dispatch few to sell more
- target influential people

Gathering the information?

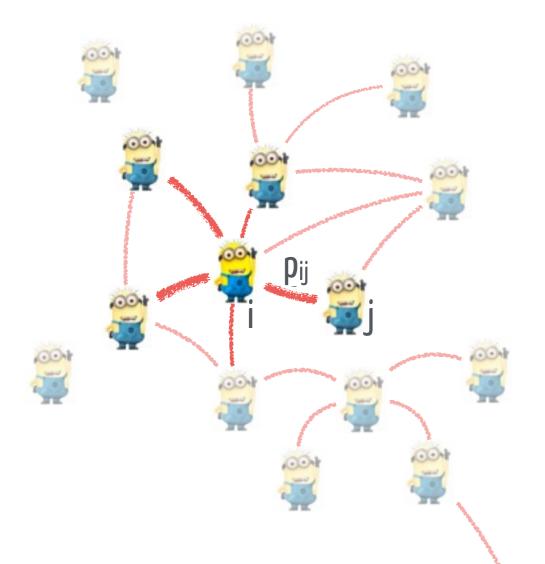
- ▶ likes on FB
- promotional codes

Unknown graphs

- all prior work needed to know the graph
- here: provably learning faster without it

REVEALING BANDITS FOR LOCAL INFLUENCE





Unknown (p_{ij})_{ij} — (symmetric) probability of influences

In each time step t = 1,, n

learner picks a node kt

environment reveals the set of influenced node Skt

Select influential people = Find the strategy maximising

$$L_n = \sum_{t=1}^n |S_{k_t,t}|$$

What this is a bandit problem?

Case n < d



PERFORMANCE CRITERION



The number of expected influences of node **k** is by definition

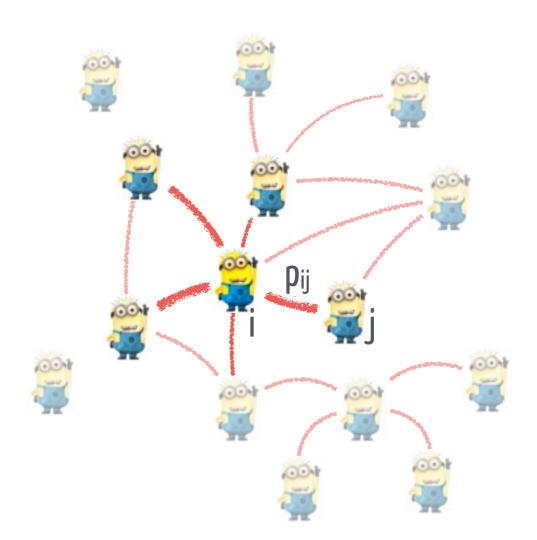
$$r_k = \mathbb{E}\left[|S_{k,t}|\right] = \sum_{j \le d} p_{k,j}$$

Oracle strategy always selects the best

$$k^* = \arg\max_{k} \mathbb{E}\left[\sum_{t=1}^{n} |S_{k,t}|\right] = \arg\max_{k} nr_k$$

Expected regret of the oracle strategy

$$\mathbb{E}\left[L_n^{\star}\right] = nr_{\star}$$



Expected regret of any adaptive strategy unaware of $(p_{ij})_{ij}$

$$\mathbb{E}\left[R_n\right] = \mathbb{E}\left[L_n^{\star}\right] - \mathbb{E}\left[L_n\right]$$

BASELINE



- We **only** receive |S| instead of S
- Can be mapped to **multi-arm** bandits
 - rewards are 0, ..., d
 - variance bounded with rkt
- We adapt **MOSS** to **GraphMOSS**
- Regret upper bound of GraphMOSS

$$\mathbb{E}\left[R_n\right] \le U \min\left(r_{\star}n, r_{\star}d + \sqrt{r_{\star}nd}\right)$$

matching lower bound



each node at least once

Crash course on stochastic bandits?

GRAPHMOSS FOR THE RESTRICTED SETTING



GraphMOSS

Input

d: the number of nodes

n: time horizon



Sample each arm twice

Update $\widehat{r}_{k,2d}$, $\widehat{\sigma}_{k,2d}$, and $T_{k,2d} \leftarrow 2$, for $\forall k \leq d$

for
$$t = 2d + 1, ..., n$$
 do

$$C_{k,t} \leftarrow 2\widehat{\sigma}_{k,t} \sqrt{\frac{\max(\log(n/(dT_{k,t})),0)}{T_{k,t}}} + \frac{2\max(\log(n/(dT_{k,t})),0)}{T_{k,t}}, \text{ for } \forall k \leq d$$

 $k_t \leftarrow \arg\max_k \widehat{r}_{k,t} + C_{k,t}$

Sample node k_t and receive $|S_{k_t,t}|$

Update $\widehat{r}_{k,t+1}$, $\widehat{\sigma}_{k,t+1}$, and $T_{k,t+1}$, for $\forall k \leq d$

end for



each node at least once



BACK TO THE REAL SETTING



- Can we actually do better?
 - Well, not really.....
 - Minimax optimal rate is still the same
- But the bad cases are somehow pathological
 - isolated nodes
 - uncorrelated being influenced and being influential
 - Barabási–Albert etc tell us that the real-world graphs are not like that
- Let's think of some measure of difficulty
 - to define some non-degenerate cases
 - ideas?

DETECTABLE DIMENSION



- number of nodes we can efficiently extract in less than n rounds
- function D controls number of nodes given a gap

$$D(\Delta) \stackrel{\text{def}}{=} |\{i \leq d : r_{\star}^{\circ} - r_{i}^{\circ} \leq \Delta\}|$$

- D(r) = d for r≥ r* and D(0) = number of most influenced nodes
- ▶ Detectable dimension $D_* = D(\Delta_*)$
- **Detectable gap** Δ * constants coming from the analysis and the Bernstein inequality

$$\Delta_{\star} \stackrel{\text{def}}{=} 16\sqrt{\frac{r_{\star}^{\circ} d \log (nd)}{T_{\star}}} + \frac{80d \log (nd)}{T_{\star}}$$

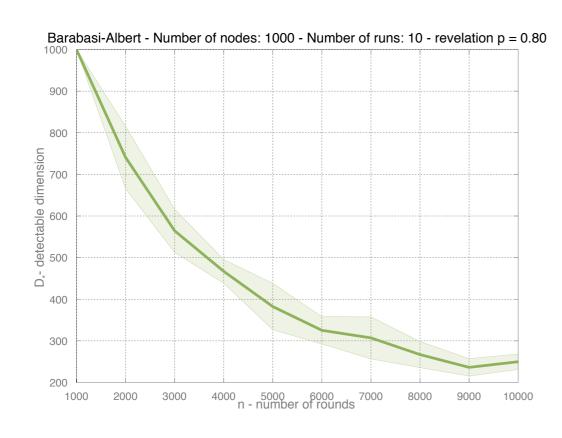
- $ilde{}$ Detectable horizon T*, smallest integer s.t. $T_{\star}r_{\star}^{\circ} \geq \sqrt{D_{\star}nr_{\star}^{\circ}}$
- Equivalently: D* corresponding to smallest T* such that

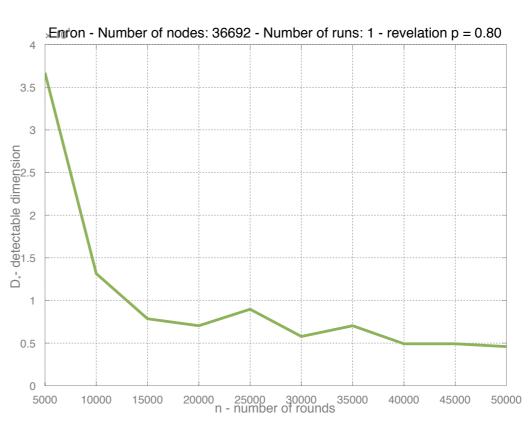
$$T_{\star}r_{\star}^{\circ} \ge \sqrt{D\left(16\sqrt{\frac{r_{\star}^{\circ}d\log\left(nd\right)}{T_{\star}}} + \frac{80d\log\left(nd\right)}{T_{\star}}\right)nr_{\star}^{\circ}}$$

HOW DOES D* BEHAVE?



- For (easy, structured) star graphs $D_* = 1$ even for small n (big gain)
- For (difficult) empty graphs D*= d even for large n (no gain)
- In general: D* roughly decreases with n and it is small when D decreases quickly
- For n large enough D* is the number of the most influences nodes
- Example: D* for Barabási–Albert model & Enron graph as a function of n





BARE SOLUTION



BAndit REvelator: 2-phase algorithm

- global exploration phase
 - super-efficient exploration 😂
 - linear regret **☞** needs to be short!
 - extracts D* nodes
- bandit phase
 - uses a minimax-optimal bandit algorithm
 - GraphMOSS is a little brother of MOSS
 - has a "square root" regret on **D**∗ nodes
- D* realizes the optimal trade-off!
 - different from exploration/exploitation tradeoff





BARE - BAndit REvelator

Input

d: the number of nodes

n: time horizon



$$T_{k,t} \leftarrow 0$$
, for $\forall k \leq d$

$$\widehat{r_{k,t}^{\circ}} \leftarrow 0$$
, for $\forall k \leq d$

$$t \leftarrow 1, \ \widehat{T}_{\star} \leftarrow 0, \ \widehat{D}_{\star,t} \leftarrow d, \ \widehat{\sigma}_{\star,1} \leftarrow d$$

Global exploration phase

while
$$t\left(\widehat{\sigma}_{\star,t} - 4\sqrt{d\log(dn)/t}\right) \leq \sqrt{\widehat{D}_{\star,t}n} \ \mathbf{do}$$

Influence a node at random (choose k_t uniformly at random) and get $S_{k_t,t}$ from this node

$$\widehat{r_{k,t+1}^{\circ}} \leftarrow \widehat{t_{t+1}} \widehat{r_{k,t}^{\circ}} + \frac{d}{t+1} S_{k_t,t}(k)$$

$$\widehat{\sigma}_{\star,t+1} \leftarrow \max_{k'} \sqrt{\widehat{r_{k',t+1}^{\circ}} + 8d \log(nd)/(t+1)}$$

$$w_{\star,t+1} \leftarrow 8\widehat{\sigma}_{\star,t+1} \sqrt{\frac{d \log(nd)}{t+1}} + \frac{24d \log(nd)}{t+1}$$

$$\widehat{D}_{\star,t+1} \leftarrow \left| \left\{ k : \max_{k'} \widehat{r_{k',t+1}^{\circ}} - \widehat{r_{k,t+1}^{\circ}} \le w_{\star,t+1} \right\} \right|$$

$$t \leftarrow t+1$$

end while

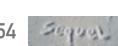
$$\widehat{T}_{\star} \leftarrow t$$
.

Bandit phase

Run minimax-optimal bandit algorithm on the $\widehat{D}_{\star,\widehat{T}_{\star}}$ chosen nodes (e.g., Algorithm 1)







EMPIRICAL RESULTS



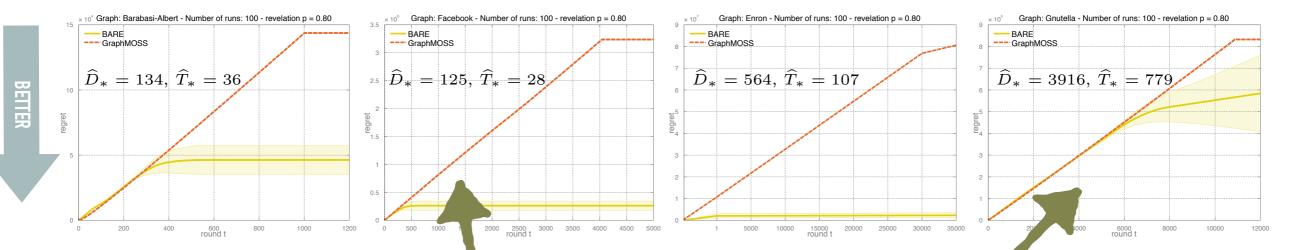
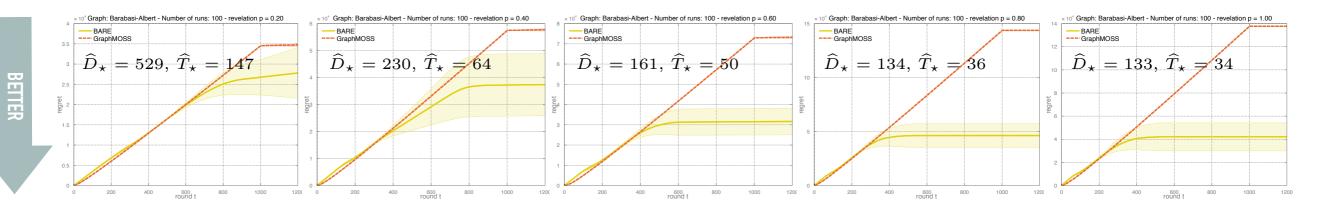


Figure 1: Left: Barabási-Albert. Middle left: Facebook. Middle right: Enren. Right: Gnutella.

Enron and Facebook vs. Gnutella (decentralised)



Varying a (constant) probability of influence

REVEALING BANDITS: WHAT DO YOU MEAN?



- Ignoring the structure again?
- $\widetilde{\mathcal{O}}\left(\sqrt{r_*TN}\right)$

reward of the

best node

- **BAndit REvelator:** 2-phase algorithm
- global exploration phase
 - super-efficient exploration
 - linear regret needs to be short!
 - extracts D* nodes
- **bandit** phase
 - uses a minimax-optimal bandit algorithm (GraphMOSS)
 - has a "square root" regret on **D*** nodes
- D* realizes the optimal trade-off!
 - different from exploration/exploitation tradeoff

Regret of BARE

$$\widetilde{\mathcal{O}}\left(\sqrt{r_*TD_*}\right)$$

- D* detectable dimension (depends on T and the structure)
 - **good case**: star-shaped graph can have D* = 1
 - **bad case:** a graph with many small cliques.
 - the worst case: all nodes are disconnected except 2

NEXT: GLOBAL INFLUENCE MODELS



- Kempe, Kleinberg, Tárdos, 2003, 2015: Independence Cascades, Linear Threshold models
 - global and multiple-source models
- Different feed-back models
 - Full bandit (only the number of influenced nodes)
 - Node-level semi-bandit (identities of influenced nodes)
 - Edge-level semi-bandit (identities of influenced edges)
 - http://arxiv.org/abs/1605.06593 (Wen, Kveton, MV)
 - IMLinUCB with linear parametrization of edge weights
 - Regret analysis for subset of graphs (forests, ...)

NEXT SEMESTER



Advanced Learning for Text and Graph Data

Time: Spring term 4 lectures and 3 Labs

Place: Polytechnique / Amphi Sauvy

Lecturer 1: Michalis Vazirgiannis (Polytechnique)

Lecturer 2: Yassine Faihe (Hewlett-Packard - Vertica)

ALTeGraD follows after Graphs in ML

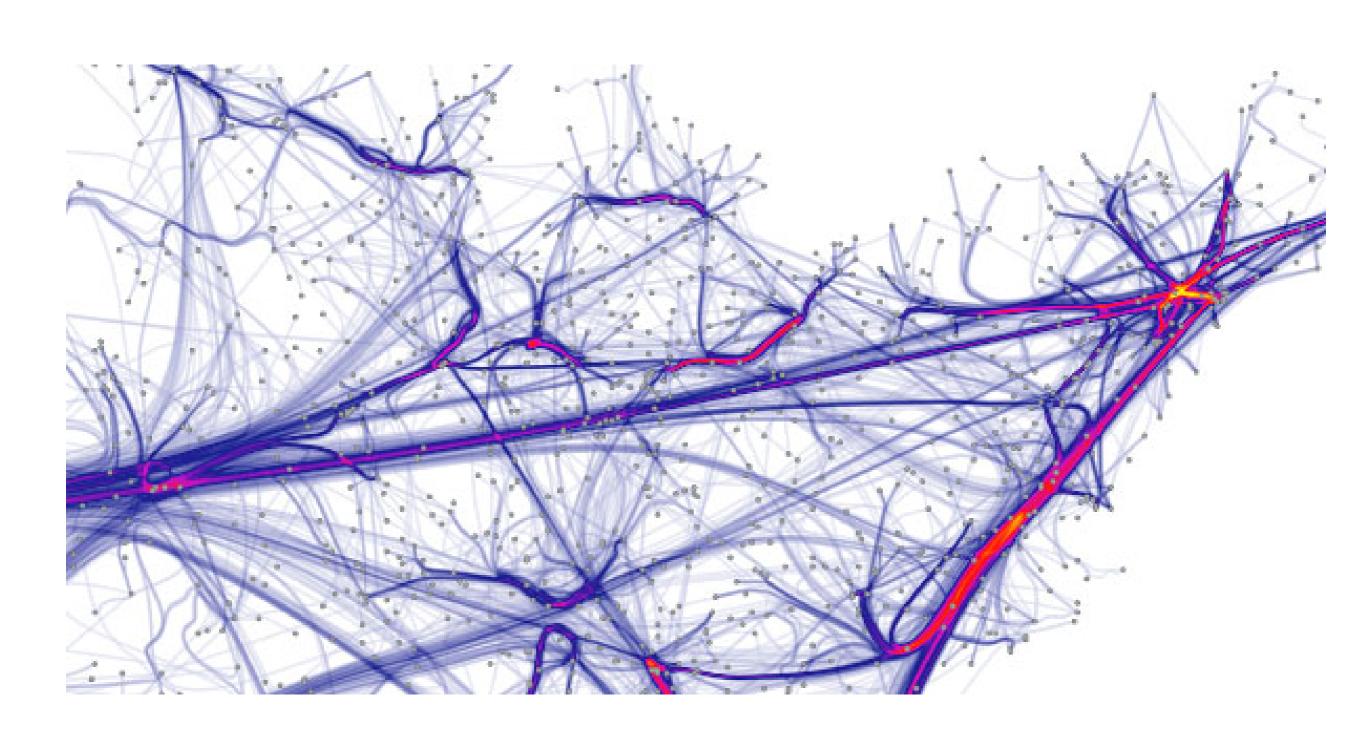
The two graph courses are coordinated to be complementary.

Some of covered graph topics not covered in this course

- Ranking algorithms and measures (Kendal Tau, NDCG)
- Advanced graph generators
- Community mining, advanced graph clustering
- Graph degeneracy (k-core & extensions)
- Privacy in graph mining

THANK YOU!





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