



# Graphs in Machine Learning

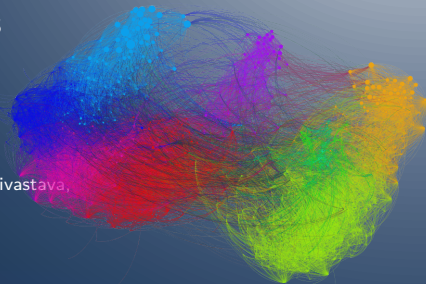
## Spectral Graph Sparsifiers in ML

Ridge Sparsifiers and Applications

Michal Valko

*Inria & ENS Paris-Saclay, MVA*

Partially based on material by: Rob Fergus, Nikhil Srivastava,  
Yiannis Koutis, Joshua Batson, Daniel Spielman



# Spectral Graph Sparsification: Intuition

What is this rescaling  $\mathbf{v}_e = \mathbf{L}_G^{-1/2} \mathbf{b}_e$  doing to the norm?

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Edges with higher  $R_{\text{eff}}$  are more **electrically significant!**

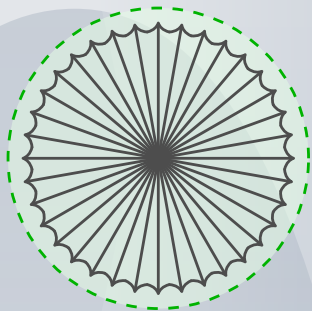
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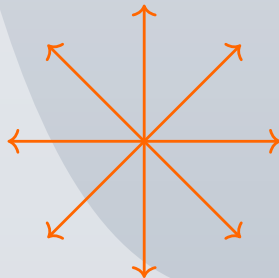


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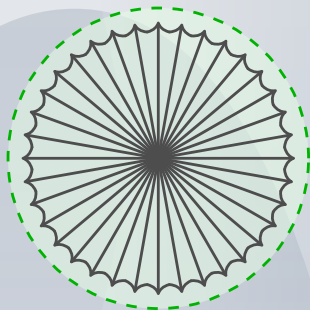
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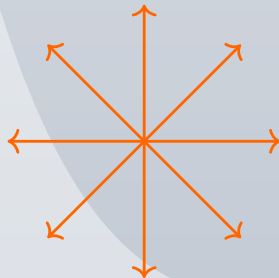
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We take a subset of these  $e_e$ s and scale them!

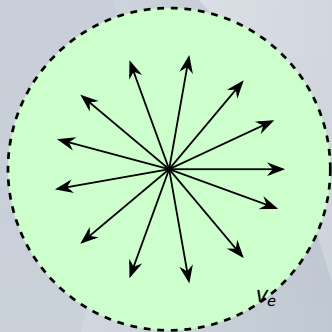
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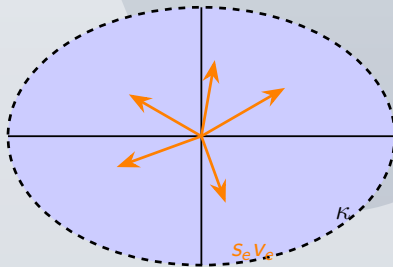
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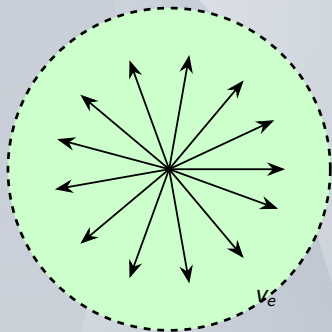
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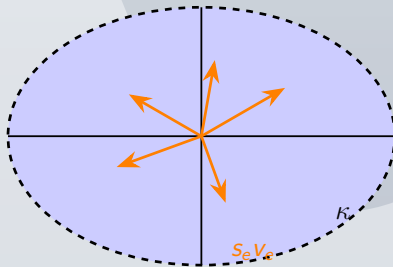
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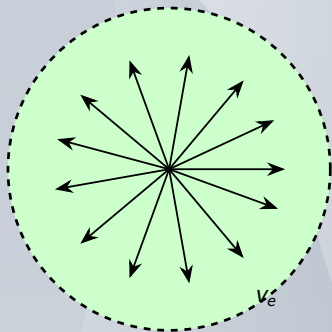


Such that the blue ellipsoid looks like identity!

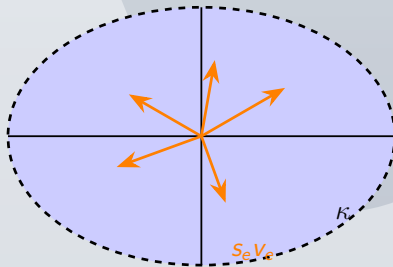
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the blue eigenvalues are between 1 and  $\kappa$

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finer bounds now available

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## Definition

**$\gamma$ -effective resistance:**  $r_e(\gamma) = \mathbf{b}_e^\top (\mathbf{L}_G + \gamma \mathbf{I})^{-1} \mathbf{b}_e$

**Effective dim.:**  $\mathbf{d}_{\text{eff}}(\gamma) = \sum_e r_e(\gamma) = \sum_{i=1}^n \frac{\lambda_i(\mathbf{L}_G)}{\lambda_i(\mathbf{L}_G) + \gamma} \leq n$

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Adapt SOA algorithm for kernel matrix approximation

SQUEAK, Calandriello et al., 2017



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↳ use sparsification internally

all the way until you hit the turtles

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↳ Use the sparsifier  $H$  to compute improved approximate  $\tilde{R}_{\text{eff}}$

Computing  $\tilde{R}_{\text{eff}}$  using the sparsifier is fast ( $m = \mathcal{O}(n \log(n))$ ), and not too many iterations are necessary.

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- Various embeddings: k-means, spectral clustering.

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`https://misovalko.github.io/mva-ml-graphs.html`

